Classwork N°.2 due to 2nd March 2012

1. Combinatory base : S and K

Proof of the completeness of the S-K basis

Question: Calculate the following combinatorial expressions using the rules:

$$\mathbf{S}xyz := xz(yz)$$

$$\mathbf{K}xy := x$$

$$\mathbf{I}_{X} := x$$

=Kx(Kx)

 $=_{\mathbf{X}}$

(2) **S (K S) K** x y z

=KSx(Kx)yz

=S(Kx)yz

=Kxz(yz)

=x(yz)

(3) **S (K (S I)) K** x y

$$=K(S I) x (K x) y$$

=SI(Kx)y

=Iy(Kxy)

 $=y(K \times y)$

=(y x)

(4) S (K (S I)) (S (K K) I) x y

2. Abstraction algorithm

T[] may be defined as follows:

$$1'$$
 $T[v] \Rightarrow v$

$$2'|$$
 $T[(E1 E2)] => (T(E1) T(E2))$

3'|
$$T[\lambda x.E] => (K T[E])$$

$$4'$$
 $T[\lambda x.x] => I$

5'|
$$T[\lambda x. \lambda y. E] \Rightarrow T[\lambda x. T[\lambda y. E]]$$

6'|
$$T[\lambda x.(E1 E2)] \Rightarrow (S T[\lambda x.E1] T[\lambda x.E2])$$

Question: Convert the lambda term $\lambda x. \lambda y. (yx)$ to a combinator:

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a)

T[λx.λy.(yx)]

=T[λx.T[λy.(yx)] (by 5)

=T[λx.(S T[λy.(y)] T[λy.x])] (by 6)

=T[λx.(S I T[λy.x])] (by 4)

=T[λx.(S I (Kx))] (by 3 and 1)

=(S T[λx.(S I)] T[λx. (Kx))] (by 6)

=(S K (S I)] T[λx. (Kx))] (by 3)

=(S K (S I)) (S T[λx. K] T[λx.x])) (by 6)
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$$=(S K (S I)) (S (K K) T[\lambda x.x])) (by 3)$$

- b) $\lambda x.\lambda y.(xy)$
- $=T[\lambda x.\lambda y.(xy)]$
- $=T[\lambda x.T[\lambda y.(xy)]]$ (by 5)
- = $T[\lambda x.(S T [\lambda y.x] T [\lambda y.y])]$ (by 6)
- $=T[\lambda x.(S T [\lambda y.x] I)]$ (by 4)
- = $K S (T[\lambda y.x] I) (by 3)$
- =K S K x I (by 3)
- =K S x (by e.-K)
- =S (by e.-K)