PA152: Efficient Use of DB 12. Replication and High Availability

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Credits

This presentation is based on:

- Materiály knihovny MSDN firmy Microsoft
- Course NoSQL databáze a Big Data management
 - Irena Holubová
 - Charles University, Prague
 - http://www.ksi.mff.cuni.cz/~holubova/NDBI040/
- PostgreSQL documentation
 - http://www.postgresql.org/docs/9.3/static/highavailability.html

Contents

- Availability
- Data distribution & Replication
- High availability
- Failover
- Recommendations

Availability





DB Server Source: Microsoft

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Determining Availability Requirements

- Hours of Operation
 - Business hours vs. all of the time
 - intranet service vs. web services
 - shift workers vs. all-around the world customers
- Connectivity Requirements
 - □ Online vs. offline applications
- Tight/Loose coupling of app and DBMS
 Synchronous vs. asynchronous data updates

Availability

Definition in operation hours

Av = "up time" / "total time" = MTTF / (MTTF+MTTR)

"up time" = the system is up and operating

□ More practical def.

Av = (total time - down time) / total time

Down time

□ Scheduled – reboot, SW/HW upgrade, ...

Unscheduled – HW/SW failure, security breaches, network unavailability, power outage, disasters, …

For "true" high-availability, down time is not distinguished

Nines

Availability as percentage of uptime Class of nines: c = [-log₁₀(1 - Av)] Assuming 24/7 operation:

Nine class	Availability	Downtime per year	Downtime per month	Downtime per week
1	90%	36.5 days	72 hours	16.8 hours
2	99%	3.65 days	7.20 hours	1.68 hours
3	99.9%	8.76 hours	43.8 minutes	10.1 minutes
4	99.99%	52.56 minutes	4.32 minutes	1.01 minutes
5	99.999%	5.26 minutes	25.9 seconds	6.05 seconds
6	99.9999%	31.5 seconds	2.59 seconds	0.605 seconds
7	99.99999%	3.15 seconds	0.259 seconds	0.0605 seconds
				Source: Wikipedia.org

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Scalability

Scalability

- Providing access to a number of concurrent users
- Handling growing amounts of data without losing performance
- □ With acceptable latency!
- Scaling Up vertical scaling → vendor dependence
 Increasing RAM
 - Multiprocessing

Scaling Out – horizontal scaling

- □ Replication
- Read-only standby servers
- Server federations / clusters / data distribution

Horizontal Scaling

- Systems are distributed across multiple machines or nodes
 - \Box Commodity machines \rightarrow cost effective
 - Often surpasses scalability of vertical approach
- Fallacies of distributed computing by Peter Deutsch
 Network
 - Is reliable, secure, homogeneous
 - Latency and transport cost is zero
 - Bandwidth is infinite
 - Topology does not change

One administrator

Source: https://blogs.oracle.com/jag/resource/Fallacies.html

Brewer's CAP Theorem

Consistency

- After an update, all readers in a distributed system see the same data
- All nodes are supposed to contain the same data at all times
- E.g. in multiple instances, all writes must be duplicated before write operation is completed.

Availability

- □ Every request receives a response
 - about whether it was successful or failed
- Partition Tolerance
 - System continues to operate despite arbitrary message loss or failure of part of the system.

Brewer's CAP Theorem

Only 2 of 3 guarantees can be given in a "shareddata" system.

Proved by Nancy Lynch in 2002

ACID

- provides Availability and Consistency
- E.g. database on a single machine

BASE

provides Availability and Partition tolerance



Source: http://bigdatanerd.wordpress.com

- Reality: you can trade a little consistency for some availability
- E.g. distributed database

BASE Properties

- Basically <u>Available</u>
 - Partial failures can occur, but without total system failure
- <u>S</u>oft state
 - System is in flux / non-deterministic
 - Changes occur all the time
- <u>Eventual consistency (replica convergence)</u>
 - □ is a liveness guarantee
 - reads eventually return the same value
 - □ is not safety guarantee
 - can return any value before it converges

Consistency

Strong (ACID) vs. Eventual (BASE) consistency

Example:

Eventual	Server A:	read(A)=1 wr	te(A,2) read(A)=2		
	Server B:	read(A)= 1	read(A)=1	read(A)=2	
	Server C:	read(A)= 1	read(A)=2		
			Inconsistent state		
	Server A:	read(A)=1 wri	rite(A,2) read(A)=2		
Strong	Server B:	read(A)= 1	read(A)=2	read(A)=2	
	Server C:	read(A)= 1	read(A)=2		

time

Replication / Distribution Model

Model of distributing data

Replication

- The same data stored in more nodes.
- □ Filtering data (sharding)
 - The data is partitioned and stored separately
 - Helps avoid replication conflicts when multiple sites are allowed to update data.

Filtering Data



Need for Distributing Data

- Brings data closer to its user
- Allows site independence
- Separates
 - Online transaction processing
 - □ Read-intensive applications
- Can reduce conflicts during user requests
 Process big data

Distribution Model

- Master-slave model (replication)
 - Load-balancing of read-intensive queries

Master node

- manages data
- distributes changes to slaves
- Slave node
 - stores data
 - queries data
 - no modifications to data



Distribution Model

- Master-master model
 - □ Typically with filtering data
 - Master for a subset of data
 - Slave for the rest

Consistency needs resolving of update conflicts



Master-master Model



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Source: Microsoft





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Replication Types

- Distributed Transactions
 - For "real" master-master model, ensures consistency
 - Low latency, high consistency
- Transactional Replication
 - Replication of incremental changes
 - □ Minimal latency

Replication Types

- Snapshot Replication
 - Periodic bulk transfer of new snapshots of data
 - Data changes substantial but infrequent
 - □ Slaves are read-only
 - □ High latency is acceptable

Replication Types

- Merge Replication
 - Autonomous changes to replicated data are later merged
 - Does not guarantee transactional consistency, but converges
 - Default and custom conflict resolution rules
 - Adv: Nodes can update data offline, sync later
 - □ Disadv: Changes to schema needed.

Maintaining High-Availability

- Standby server
 - □ Shared disk failover (NAS)
 - □ File system replication (DRBD)
 - Transaction log shipping
 - Trigger-based replication
 - Statement-Based
 Replication
 Middleware



Log-shipping Standby Server

- Also called warm standby
- Primary node
 - serves all queries
 - in permanent archiving mode
 - Continuous sending of WAL records to standby servers
- Standby server
 - serves no queries
 - in permanent recovery mode
 - Continuous processing of WAL records arriving from primary node
- Log shipping can be synchronous/asynchronous
- Disadvantage: all tables are replicated typically
- Advantage: no schema changes, no trigger definitions

Failover

- If primary fails, standby server begins failover.
 - Standby applies all WAL records pending,
 - □ marks itself as primary,
 - □ starts to serve all queries.
- If standby fails, no action taken.
 - After becoming online, catch-up procedure is started.
- Heartbeat mechanism
 - □ to continually verify the connectivity between the two and the viability of the primary server

Failover

Failover by standby succeeded New standby should be configured

Original primary node becomes available

- $\Box \rightarrow$ inform it that it is no longer the primary
 - do so-called STONITH (Shoot The Other Node In The Head),
 - otherwise serious data corruption/loss may occur

Typically old primary becomes new standby

Primary and Standby Servers

- Swap primary and standby regularly
 - □ To verify recovery steps
 - To do necessary maintenance on standby server
 - SW/HW upgrades, ...

Recommended Practices

- Maximize availability at each tier of the application
- Keep standby servers on a different subnet
- Different power supply to the primary server
- Test whether your availability solution works

NewSQL

Distributed database that scales out

CP system

- trades availability for consistency when partition happens
- MySQL cluster, Google Spanner, VoltDB, ...
 In fact, master-master replication with data sharding