SOLUTIONS

Exercises on Block2:

Finding Frequent Item Sets
Finding Similar Items
Searching in Data Streams

Advanced Search Techniques for Large Scale Data Analytics

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Frequent Item Sets (1) – Assignment

- Suppose 100 items (numbered 1 to 100) and 100 baskets (numbered 1 to 100)
  - Item *i* is in basket *b* if and only if *i* divides *b* with no remainder, i.e., item 1 is in all the baskets, item 2 is in all fifty of the even-numbered baskets, etc.
- Consider that the support threshold is 5:
  1) Identify the frequent items
  2) Compute the confidence of these association rules
     a) $\{5, 7\} \rightarrow 2$
     b) $\{2, 3, 4\} \rightarrow 5$

Frequent Item Sets (1) – Recap

- **Simplest question:** Find sets of items that appear together "frequently" in baskets
- ***Support*** for item set I: Number of baskets containing all items in I
  - (Often expressed as a fraction of the total number of baskets)
- Given a ***support threshold*** s, then sets of items that appear in at least s baskets are called ***frequent itemsets***

TID	*Items*
1	Bread, Coke, Milk
2	Beer, Bread
3	Beer, Coke, Diaper, Milk
4	Beer, Bread, Diaper, Milk
5	Coke, Diaper, Milk

Support of {Beer, Bread} = 2

Frequent Item Sets (1) – Recap

- **Association Rules:**
  If-then rules about the contents of baskets
- $\{i_1, i_2, \ldots, i_k\} \rightarrow j$ means: "if a basket contains all of $i_1, \ldots, i_k$ then it is ***likely*** to contain j"
- **In practice there are many rules, want to find significant/interesting ones!**
- ***Confidence*** of this association rule is the probability of j given $I = \{i_1, \ldots, i_k\}$

$$\operatorname{conf}(I \rightarrow j) = \frac{\operatorname{support}(I \cup j)}{\operatorname{support}(I)}$$

Frequent Item Sets (1) – Solution

1) 20 frequent items: **1–20**
2) Association rules:
   a) The baskets containing both items 5 and 7 are baskets 35 and 70, in which only basket 70 also contains item 2. Hence, the confidence of the rule {5, 7} → 2 is **1/2**.
   b) The baskets whose numbers are the multiples of 12 contain item set {2, 3, 4} as a subset – there are 8 such baskets. The baskets whose numbers are the multiples of 60 contain item set {2, 3, 4, 5} as a subset – there is 1 such basket. Hence, the confidence of the rule {2, 3, 4} → 5 is **1/8**.

Frequent Item Sets (2) – Assignment

- Consider the following twelve baskets, each of them contains 3 of 6 items (1 through 6):
  - {1, 2, 3} {2, 3, 4} {3, 4, 5} {4, 5, 6}
  - {1, 3, 5} {2, 4, 6} {1, 3, 4} {2, 4, 5}
  - {3, 5, 6} {1, 2, 4} {2, 3, 5} {3, 4, 6}
- Suppose the support threshold is 4. On the first pass of the PCY algorithm, a hash table with 11 buckets is used, and the set $\{i, j\}$ is hashed to bucket $i \times j \bmod 11$:
  1) Compute the support for each item and each pair of items
  2) Which pairs hash to which buckets?
  3) Which buckets are frequent?
  4) Which pairs are counted on the second pass?

Frequent Item Sets (2) – Recap

PCY Algorithm – First Pass

New in PCY

```
FOR (each basket) :
    FOR (each item in the basket) :
         add 1 to item's count;
    FOR (each pair of items) :
         hash the pair to a bucket;
         add 1 to the count for that bucket;
```

- **Few things to note:**
  - Pairs of items need to be generated from the input file; they are not present in the file
  - We are not just interested in the presence of a pair, but we need to see whether it is present at least ***s*** (support) times

Frequent Item Sets (2) – Recap

- **Observation: If a bucket contains a frequent pair, then the bucket is surely frequent**
- However, even without any frequent pair, a bucket can still be frequent ☹
  - So, we cannot use the hash to eliminate any member (pair) of a “frequent” bucket
- **But, for a bucket with total count less than s, none of its pairs can be frequent ☺**
  - Pairs that hash to this bucket can be eliminated as candidates (even if the pair consists of 2 frequent items)
- **Pass 2:** Only count pairs that hash to frequent buckets

Frequent Item Sets (2) – Solution 1/4

1) Compute the support for each item and each pair of items

- Support for each item:

item	1	2	3	4	5	6
support	4	6	8	8	6	4

- Support for each pair of items:

pair	{1, 2}	{1, 3}	{1, 4}	{1, 5}	{1, 6}	{2, 3}	{2, 4}	{2, 5}
support	2	3	2	1	0	3	4	2

pair	{2, 6}	{3, 4}	{3, 5}	{3, 6}	{4, 5}	{4, 6}	{5, 6}
support	1	4	4	2	3	3	2

Frequent Item Sets (2) – Solution 2/4

2) Which pairs hash to which buckets?

- The set {*i*, *j*} is hashed to bucket no.: $i \times j$ mod 11

pair	{1, 2}	{1, 3}	{1, 4}	{1, 5}	{1, 6}	{2, 3}	{2, 4}	{2, 5}
bucket	2	3	4	5	6	6	8	10

pair	{2, 6}	{3, 4}	{3, 5}	{3, 6}	{4, 5}	{4, 6}	{5, 6}
bucket	1	1	4	7	9	2	8

Frequent Item Sets (2) – Solution 3/4

3) Which buckets are frequent?

- Bucket support – sum of supports of pairs belonging to the given bucket:

bucket	0	1	2	3	4	5	6	7
support	0	5	5	3	6	1	3	2

bucket	8	9	10
support	6	3	2

- The frequent buckets are those with support above 4, i.e., buckets: 1, 2, 4, 8

Frequent Item Sets (2) – Solution 4/4

4) Which pairs are counted on the second pass
   - As only pairs in frequent buckets will be counted on the second pass of PCY, they are:

     {1, 2}, {1, 4}, {2, 4}, {2, 6}, {3, 4}, {3, 5}, {4, 6}, {5, 6}

Shingling (1) – Assignment

- Consider two documents A and B
  - If their 3-shingle resemblance is 1 (using Jaccard similarity), does that mean that A and B are identical?
    - If so, prove it. If not, give a counterexample.

Shingling (1) – Recap

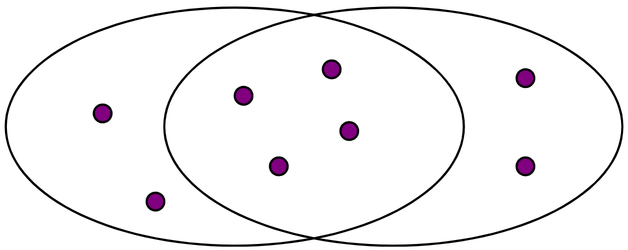
- A *k*-shingle (or *k*-gram) for a document is a sequence of *k* tokens that appears in the doc
  - Tokens can be characters, words or something else, depending on the application
  - Assume tokens = characters for examples

- **Example:** **k=2**; document $\mathbf{D_1}$ = abcab
  Set of 2-shingles: $\mathbf{S(D_1)}$ = {ab, bc, ca}
  - **Option:** Shingles as a bag (multiset), count ab twice: $\mathbf{S'(D_1)}$ **=** {ab, bc, ca, ab}

Shingling (1) – Recap

- **Document D_1 is a set of its k-shingles C_1=S(D_1)**
- Equivalently, each document is a 0/1 vector in the space of *k*-shingles
  - Each unique shingle is a dimension
  - Vectors are very sparse
- **A natural similarity measure is the Jaccard similarity:**

$$sim(D_1, D_2) = |C_1 \cap C_2| / |C_1 \cup C_2|$$



Shingling (1) – Solution

- No, the documents A and B need not be identical
  - Counterexample:
    - A: **abab**
      - 3-shingles: S(A) = {aba, bab}
    - B: **baba**
      - 3-shingles: S(B) = {bab, aba}
      - $sim(A, B) = | S(A) \cap S(B) | / | S(A) \cup S(B) | = 1$

Shingling (2) – Assignment

- Consider two documents A and B
  - Each document's number of token is $O(n)$
    - It does not matter whether tokens are characters or words
  - What is the runtime complexity of computing A and B's k-shingle resemblance (using Jaccard similarity)?
    - Assume that comparison of two k-shingles to assess their equivalence is $O(k)$
    - Express your answer in terms of n and k, where $n >> k$

Shingling (2) – Solution

- Time to create shingles: $O(n)$
- Time to find intersection (using the brute force algorithm): $O(k \cdot n^2)$
  - n shingles in each document
- Time to find union (using the intersection): $O(n)$
- Total time: $\boldsymbol{O(k \cdot n^2)}$

Finding Similar Items (1) – Assignment

- Compute the Jaccard similarities of each pair of the following three sets:
  - A = {1, 2, 3, 4}
  - B = {2, 3, 5, 7}
  - C = {2, 4, 6}

Finding Similar Items (1) – Solution

- *sim*(A, B) = 2/6 = 1/3
- *sim*(A, C) = 2/5
- *sim*(B, C) = 1/6

Finding Similar Items (2) – Assignment

- For the matrix

Element	D1	D2	D3	D4
0	0	1	0	1
1	0	1	0	0
2	1	0	0	1
3	0	0	1	0
4	0	0	1	1
5	1	0	0	0

1) Compute the minhash signature for each column (document) using the following hash functions:
   - $h_1(x) = 2x + 1 \bmod 6$
   - $h_2(x) = 3x + 2 \bmod 6$
   - $h_3(x) = 5x + 2 \bmod 6$
2) Which of these hash functions are true permutations?
3) How close are the estimated Jaccard similarities for the six pairs of columns to the true Jaccard similarities?

Finding Similar Items (2) – Recap

- **Rows** = elements (e.g., shingles)
- **Columns** = sets (e.g., documents)
  - 1 in row ***e*** (shingle) and column ***s*** (document) if and only if ***e*** is a member of ***s***
  - Column similarity is the Jaccard similarity of the corresponding sets (rows with value 1)
  - **Typical matrix is sparse!**
- **Each document is a column:**
  - **Example: sim(C_1 ,C_2) = ?**
    - Size of intersection = 3; size of union = 6, Jaccard similarity (not distance) = 3/6
    - **d(C_1,C_2) = 1 – (Jaccard similarity) = 3/6**

Documents

Shingles

1	1	1	0
1	1	0	1
0	1	0	1
0	0	0	1
1	0	0	1
1	1	1	0
1	0	1	0

Finding Similar Items (2) – Recap

Min-Hashing Example

2nd element of the permutation is the first to map to a 1

Permutation π

2	4	3
3	2	4
7	1	7
6	3	2
1	6	6
5	7	1
4	5	5

Input matrix (Shingles x Documents)

1	0	1	0
1	0	0	1
0	1	0	1
0	1	0	1
0	1	0	1
1	0	1	0
1	0	1	0

Signature matrix *M*

2	1	2	1
2	1	4	1
1	2	1	2

4th element of the permutation is the first to map to a 1

Finding Similar Items (2) – Solution 1+2/3

1) Compute the minhash signature for each column using the following hash functions:
   - $h_1(x) = 2x + 1 \bmod 6$
   - $h_2(x) = 3x + 2 \bmod 6$
   - $h_3(x) = 5x + 2 \bmod 6$

   Hashes are computed on element IDs:

Element	D1	D2	D3	D4	$h_1(x)$	$h_2(x)$	$h_3(x)$
0	0	1	0	1	**1**	**2**	**2**
1	0	1	0	0	**3**	**5**	**1**
2	1	0	0	1	**5**	**2**	**0**
3	0	0	1	0	**1**	**5**	**5**
4	0	0	1	1	**3**	**2**	**4**
5	1	0	0	0	**5**	**5**	**3**

Minhash signature:

D1	D2	D3	D4
5	1	1	1
2	2	2	2
0	1	4	0

(rows correspond to hash functions)

2) Which of these hash functions are true permutations: $\boldsymbol{h_3}$ only

Finding Similar Items (2) – Solution 3/3

3) How close are the estimated Jaccard similarities for the six pairs of columns (documents) to the true Jaccard similarities?

Jaccard similarities on	D1 / D2	D1 / D3	D1 / D4	D2 / D3	D2 / D4	D3 / D4
Original documents	0	0	0.25	0	0.25	0.25
Minhash signatures	0.33	0.33	0.67	0.67	0.67	0.67

- => the estimated Jaccard similarities are not close to the true ones at all
  - To make the estimated similarity closer to the true one, there is a need of more and better (i.e., resulting in true permutations) hash functions

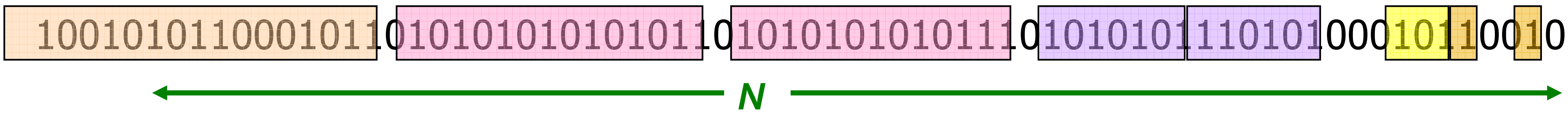
Data Streams (1) – Assignment

- Suppose we are maintaining a count of 1s using the DGIM method
  - Each bucket is represented by (*i*, *t*)
    - *i* – the number of 1s in the bucket
    - *t* – the bucket timestamp (time of the most recent 1)
- Consider the following properties:
  - Current time is 200
  - Window size is 60
  - Current buckets are:
    - (16, 148) (8, 162) (8, 177) (4, 183) (2, 192) (1, 197) (1, 200)
  - At the next ten clocks (201 through 210), the stream has 0101010101
- What will the sequence of buckets be at the end of these ten inputs?

Data Streams (1) – Recap

DGIM: Buckets

- A ***bucket*** in the DGIM method is a record consisting of:
  - **(A) The timestamp of its end** [O(log *N*) bits]
  - **(B) The number of 1s between its beginning and end** [O(log log *N*) bits]
- **Constraint on buckets:**
  Number of **1s** must be a power of **2**
  - That explains the **O(log log *N*)** in **(B) above**



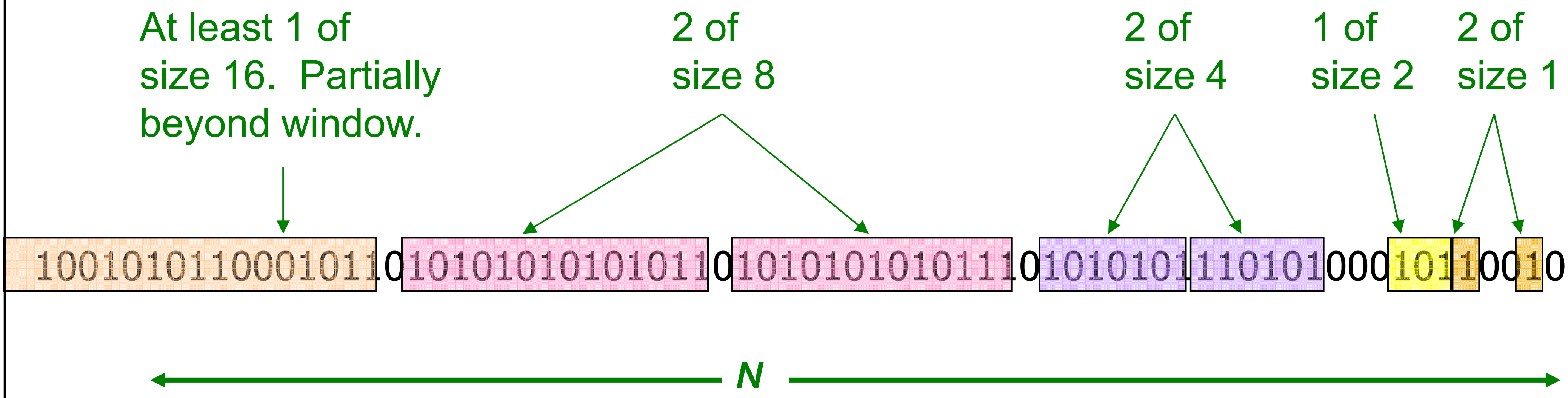
Data Streams (1) – Recap

Representing a Stream by Buckets

- Either **one** or **two** buckets with the same **power-of-2 number** of **1s**
- **Buckets do not overlap in timestamps**
- **Buckets are sorted by size**
  - Earlier buckets are not smaller than later buckets
- Buckets disappear when their end-time is **>** ***N*** time units in the past

Data Streams (1) – Recap

Example: Bucketized Stream



Three properties of buckets that are maintained:

- Either one or two buckets with the same **power-of-2** number of **1s**
- Buckets do not overlap in timestamps
- Buckets are sorted by size

Data Streams (1) – Recap

Updating Buckets (1)

- When a new bit comes in, drop the last (oldest) bucket if its end-time is prior to ***N*** time units before the current time
- **2 cases:** Current bit is **0** or **1**
- **If the current bit is 0:**
  no other changes are needed

Data Streams (1) – Recap

Updating Buckets (2)

- **If the current bit is 1:**
  - **(1)** Create a new bucket of size **1**, for just this bit
    - **End timestamp = current time**
  - **(2)** If there are now **three buckets of size 1**, **combine the oldest two into a bucket of size 2**
  - **(3)** If there are now **three buckets of size 2**, **combine the oldest two into a bucket of size 4**
  - **(4) And so on …**

Data Streams (1) – Recap

Updating buckets (example):

Current state of the stream:

1001010110001011010101010101010110101010101010111010101011101010001011001 0

Bit of value 1 arrives

0010101100010110101010101010101101010101010101110101010111010100010110010 1

Two orange buckets get merged into a yellow bucket

0010101100010110101010101010101101010101010101110101010111010100010110010 1

Next bit 1 arrives, new orange bucket is created, then 0 comes, then 1:

01011000101101010101010101010110101010101010111010101011101010001011001011 01

Buckets get merged…

01011000101101010101010101010110101010101010111010101011101010001011001011 01

State of the buckets after merging

01011000101101010101010101010111010101010101011101010101110101000101100101101

Data Streams (1) – Solution

- There are 5 occurrences of 1s in the stream. Each one updates the buckets to be:
  - (1) Combine the oldest two buckets of size 1

    (16, 148) (8, 162) (8, 177) (4, 183) (2, 192) (1, 197) (1, 200) (1, 202)

    => (16, 148) (8, 162) (8, 177) (4, 183) (2, 192) (2, 200) (1, 202)
  - (2) No combination needed

    (16, 148) (8, 162) (8, 177) (4, 183) (2, 192) (2, 200) (1, 202) (1, 204)
  - (3) Combine the oldest two buckets of size 1, and then oldest two buckets of size 2

    (16, 148) (8, 162) (8, 177) (4, 183) (2, 192) (2, 200) (1, 202) (1, 204) (1, 206)

    => (16, 148) (8, 162) (8, 177) (4, 183) (2, 192) (2, 200) (2, 204) (1, 206)

    => (16, 148) (8, 162) (8, 177) (4, 183) (4, 200) (2, 204) (1, 206)
  - (4) No combination needed; window size is 60, so (16, 148) should be dropped

    (16, 148) (8, 162) (8, 177) (4, 183) (4, 200) (2, 204) (1, 206) (1, 208)

    => (8, 162) (8, 177) (4, 183) (4, 200) (2, 204) (1, 206) (1, 208)
  - (5) Combine the oldest two buckets of size 1

    (8, 162) (8, 177) (4, 183) (4, 200) (2, 204) (1, 206) (1, 208) (1, 210)

    => **(8, 162) (8, 177) (4, 183) (4, 200) (2, 204) (2, 208) (1, 210)**