IA038 Types and Proofs

1. History of Math & Motivations

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February 21, 2023

Nature of Mathematical Objects

- Plato's "Academy": "*Let no one who is not a geometer enter*" Plato (427-348 BCE)
- Platonism:
- 1. There are mathematical objects
- 2. These are abstract objects existing outside of space and time
- 3. Math objects always existed & are entirely independent od people
- Math objects do not interact with the physical world in any "causal" way – we cannot change them, and they cannot change us, yet
- 5. We are somehow able to gain knowledge of them

Prehistory of Formal Reasoning

- Aristotle (384-322 BCE) Analytica Posteriora as a deductive science from basic truths or axioms
 * Deductive proofs as demonstration arguments in Euclid geometry
 * Logic in the form of syllogisms independent of mathematical/geometrical proofs
- Sufficient for more than two millenia

Logicism

- Gottfried Wilhelm Leibniz (1646-1716) Mathematical facts as truths of reason Hoped in "calculus ratiocinator" as a systematic calculational logic for representation of human reasoning (similar to the differential calculus for mathematical physics) Introduced logical operations (but did not publish this)
- Immanuel Kant (1724-1804) Critique of Pure Reason Arithmetic and Geometry as synthetic a priori issues akin to Metaphysics
- Richard J. W. Dedekind (1831-1916) culmination of arithmetization of Arithmetic and Geometry

Logicism

- Gottlob Frege (1848-1925)
- Begriffsschrift, 1879 (concept notation for pure thought / logic)
- Grundlagen der Arithmetic, 1884
- \succ Analyticity of arithmetic truths derived from their
- > 7+5=12 as analytical truth (contrary to Kant)
- Grundsetze der Arithmetic, Vol. 1, 1893
- \succ distinction between sense and reference
- \succ introduction of notation for concepts and semantics
- symbolic language for expressing everything explicitly & finite set of rules
- Vol. 2 was at the publisher in 1903, when Russel wrote to Frege about the Russel Paradox (A={B: B is a set & B∉B}, and both A∉A, and A∉A)

Problems with infinite sets in logicism, also Gödel's incompleteness

Formalization of logical inference

Giuseppe Peano (1858-1932)
Around 1890 formalization of logical inference
Formal rules based on axioms and Modus
Ponens (A=>B, A |- B)

• Bertrand Russel (1872-1970) Principia Mathematica, 1910-13, with Whitehead: expressing axioms as basic truths, and deriving logical truths by Modus Ponens and universal generalization David Hilbert (1862-1943) Grundlagen der Geometrie,1899

Four foundational problems:

- 1. Formalization of mathematical theory
- 2. Proof of consistency of the the axioms
- 3. Independence and completeness of the axioms
- 4. The decision problem: is there a method answering any question in the theory?

David Hilbert

- By 1920's, "Hilbert style" axiomatic approach dominates
- Propositional logic proved complete and decidable
- Predicate logic presented in Hilbert style by Ackermann by 1920

Hilbert Program ~1920

- Hilbert Program: expressing higher mathematics in terms of elementary Arithmetics; formalizing all Mathematics in axiomatic form together with a proof of completeness (finitistic methods, purely intuitive basis of concrete signs)
- P. Bernays, W. Ackermann, J. von Neumann, J. Herbrand
- Ackermann and von Neumann proof of consistency of number theory, Ackermann thought near completion for analysis
- Hilbert claimed in 1928 in Bologna that the work is essentially completed

Kurt Gödel: completeness of First-Order Predicate Logic

- Completeness of First-Order Predicate Logic stated by Hilbert and Ackermann in 1928
- Kurt Gödel tackled this in his doctoral thesis in 1929
- Thm: Every logical expression is either satisfiable or refutable, aka Every valid logical expression is provable
- Presented as his forthcoming PhD thesis in September 1930 in Königsberg

Kurt Gödel: incompleteness of formal systems

- Also in September 1930 in Königsberg, presented as an "aside" (not a talk on the conference programme)
- The First Incompleteness Thm showing existence of arithmetic formulas neither povable nor refutable in Peano arithmetic
- The Second Incompleteness Thm showing that consistency of arithmetic cannot be proved in Arithmetics itself, $Con(P) \equiv \neg Prov([0]=[1])$
- von Neumann interrupted his lectures on Hilbert proof theory in the Fall of 1930; seeing Hilbert program could not be achieved at all

 \rightarrow Destruction of Hilbert Program (impossibility of proving consistence of a formal system inside of it)

Intuitionism (constructivism)

• L.E.J. Brower (1881-1966): mathematical knowledge comes from constructing mathematical objects within human intuition; belief that the law of wxcluded middle, or indirect existential proofs are dangerous to the coherence of Mathematics

Around 1908 – clash with Hilbert, also in "The untrustworthiness of the principles of logic" challenged the belief that the rules of classical logic which came from Aristotle have absolute validity

• Arend Heyting (1898-1980) was his PhD student developing intuitionism further

Gödel and Gentzen

• Translation from Peano arithmetic to intuitionistic Heyting arithmetic by Gödel, in parallel with Gentzen, 1932-33

Gentzen

- Gentzen thesis (1934-35) on analysis of mathematical proofs
- Natural Deduction (intro & elim rules, esp. suitable for intuitionistic logic)

$$\frac{A \ B}{A \ \& B} \& \mathbf{I} \qquad \qquad \frac{A \ \& B}{A} \& \mathbf{E}_1 \qquad \frac{A \ \& B}{B} \& \mathbf{E}_2$$

$$\begin{bmatrix} A \\ \vdots \\ \\ \frac{B}{A \supset B} \supset I \end{bmatrix}$$

$$\frac{A \supset B \quad A}{B} \supset \mathbf{E}$$

Gentzen

- Gentzen thesis (1934-35) on analysis of mathematical proofs
- "Sequenzenkalkul", Sequent Calculus
- Normalization and cut-elimination

$$\frac{\Gamma \to \Delta, A \quad \Gamma \to \Delta, B}{\Gamma \to \Delta, A \& B} \operatorname{R\&} \frac{A, \Gamma \to \Delta}{A \& B, \Gamma \to \Delta} \operatorname{L\&}_1 \frac{B, \Gamma \to \Delta}{A \& B, \Gamma \to \Delta} \operatorname{L\&}_2$$

Implication

$$\frac{A, \Gamma \rightarrow \Delta, B}{\Gamma \rightarrow \Delta, A \supset B} R \supset$$

$$\frac{\Gamma \rightarrow \Theta, A \quad B, \Delta \rightarrow \Lambda}{A \supset B, \Gamma, \Delta \rightarrow \Theta, \Lambda} L \supset$$

$$\frac{\Gamma \to A \quad A, \Delta \to C}{\Gamma, \Delta \to C} \operatorname{Cut}$$

Gödel and Gentzen

- Gentzen gave an alternative proof of the Incompleteness Thm (written 1939, published 1943) by showing a formula unprovable in Peano arithmetic (thus also showing consistency of Peano arithmetic)
- Gödel's proof of consistency using Dialectica interpretation

(Unclear mutual interaction in 1939.)

Computability and Undecidable problems

- Hilbert (1928): "Entscheidungsproblem": Is there a general effective procedure deciding whether or not a given formula A of a calculus is provable?
- 1936: Alonzo Church proved on the basis of λ-calculus and Alan Turing on the basis of Turing machines several months later that

the answer to the Entscheidungsproblem is negative

• Existence of *undecidable problems* in Informatics

Computability and Undecidable problems

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- Existence of *undecidable problems* in Informatics
- Church-Turing thesis (Kleene, 1952) computable functions / computability

Alonzo Church (1903–1995) and his λ-calculus

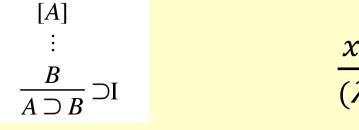
- λ -abstraction: making bound variables in function definitions explicit
- λx(M) means definition of a function mapping an argument a into M[x/a] ("formal parameters" in programming languages)
- Operational semantics via rewriting: $\lambda x(M)N \rightarrow M[x/N]$
- Fixed-point Theorem: For each F there is X s.t. FX= X
- There is a fixed-point combitator \mathbf{Y} s.t. $F(\mathbf{Y}F) = \mathbf{Y}F$
- Proof: $\mathbf{Y} = \lambda f(\lambda x.f(xx))(\lambda x.f(xx))$

Church, Turing, and Gödel

- Church using λ -calculus as a formal tool tried to formalize mathematics
- Learning about Gödel's result, claimed that it does not apply to this system
- Kleene recursive functions
- Rosser reductions
- Church-Turing thesis (Kleene, 1952) for computable functions / computability:
 - Turing machines
 - λ -definability
 - Gödel's general recursive functions (Princeton, 1934)

Curry-Howard correspondence aka "formulae-as-types"

• Computational semantics for intuitionistic logic



$$\frac{x^A \quad M^B}{(\lambda x.M)^{A \to B}}$$

$$\frac{A \supset B \quad A}{B} \supset \mathbf{E}$$

$$\frac{M^{A \to B} \quad N^A}{(MN)^B}$$

Curry-Howard correspondence aka "formulae-as-types"

- Computational semantics for intuitionistic logic
- Computations = normalization

$$\begin{bmatrix} A \\ \vdots \\ B \\ A \supset B \\ B \\ B \\ \end{bmatrix} \longrightarrow \begin{bmatrix} A \\ \vdots \\ B \\ \end{bmatrix} \xrightarrow{A \land MB} N^{A} \longrightarrow M[x/N]$$

Curry-Howard correspondence aka "formulae-as-types"

- Computational semantics for intuitionistic logic
- Computations = normalization
- Intuitionistic logic not tied to any philosophy of Mathematics, but corresponds to program execution
- Girard's Linear logic as a Sequent-Calculusstyle system capable expressing parallel operations (via proof normalization)

Completing Gödel's rupture in the structure of mathematics

- Gödel opened a crack in the foundations of Mathematics
- The complement of this rupture lays outside of combinatorial formulation; still within the real of Mathematics
- The "inside" of this crack opened up a new discipline, Informatics; may be thought of as a camouflage of formal logic (proof theory) into fairly applicable computational tool

Proof theory as a basis for constructivistic formulation

- Frege vs. Hilbert concerning Platonism vs. Formalism (Hilbert Mathematics is invented and best viewed as formal symbolic games without intrinsic meaning)
- Hilbert vs. Brower concerning Formalism vs. Intuitionism
- Constructivism (Kronecker (1823-1891) with "God made the natural numbers, all else is the work of man", and esp. Andrej Markov who claimed Mathematics should deal exclusively with constructive objects)
- Proof generation as object construction
- Proof simplification/normalization as a computational mechanism (even without underlying formula semantics)

Textbook for this course

Proofs and Types

Jean-Yves Girard, Yves Lafont and Paul Taylor

Cambridge University Press (Cambridge Tracts in Theoretical Computer Science, 7), ISBN <u>0 521 37181 3</u>; first published 1989, reprinted with corrections 1990

PDF available from http://www.paultaylor.eu/stable/Proofs+Types.html

- 1. Sense, Denotation and Semantics
- 2. Natural Deduction
- 3. The Curry-Howard Isomorphism
- 4. The Normalisation Theorem
- 5. Sequent Calculus
- 6. Strong Normalisation Theorem
- 7. Gödel's system T
- 8. Coherence Spaces
- 9. Denotational Semantics of T
- 10. Sums in Natural Deduction
- 11. System F
- 12. Coherence Semantics of the Sum
- 13. Cut Elimination (Hauptsatz)
- 14. Strong Normalisation for F
- 15. Representation Theorem

Completion Requirements

- Essay on a topic using this concept, developing necessary mathematical details.
- Structure of the essay corresponding to an article introducing the topic, and developing details.
- ➤ This may either be something relevant to your work/interest, or e.g. taking some of the results concerning normalization within a suitable formal system, and completing proofs in sufficient mathematical detail.
- > 3-5 thousand words (approx. 6-12 pages).