

Untyped lambda calculus

terms and values

$$M ::= x \mid MM' \mid \lambda x.M$$
$$V ::= \lambda x.M$$

full β -reduction evaluation rules

$$\frac{}{(\lambda x.M)N \rightarrow_{\beta} M[x := N]} \quad \frac{M_1 \rightarrow_{\beta} M_2}{\lambda x.M_1 \rightarrow_{\beta} \lambda x.M_2}$$
$$\frac{M_1 \rightarrow_{\beta} M_2}{M_1 N \rightarrow_{\beta} M_2 N} \quad \frac{N_1 \rightarrow_{\beta} N_2}{M N_1 \rightarrow_{\beta} M N_2}$$

call-by-value evaluation rules

$$\frac{}{(\lambda x.M)V \rightarrow_{\beta} M[x := V]} \quad \frac{M_1 \rightarrow_{\beta} M_2}{M_1 N \rightarrow_{\beta} M_2 N} \quad \frac{N_1 \rightarrow_{\beta} N_2}{V N_1 \rightarrow_{\beta} V N_2}$$

Church numerals and booleans

$$\underline{0} := \lambda f.\lambda x.x \qquad \text{true} := \lambda x.\lambda y.x$$
$$\underline{1} := \lambda f.\lambda x.f x \qquad \text{false} := \lambda x.\lambda y.y$$
$$\underline{n} := \underbrace{f(f \dots (f(x) \dots))}_{n\text{-times}} = \lambda f.\lambda x.f^n(x)$$

Simply-typed lambda calculus λ^{\rightarrow} with Booleans

terms and values

$$t ::= x \mid t t' \mid \lambda x : T.t \mid \text{true} \mid \text{false} \mid \text{if } t \text{ then } t' \text{ else } t''$$
$$v ::= \lambda x : T.t \mid \text{true} \mid \text{false}$$

types

$$T ::= T \rightarrow T \mid \text{Bool}$$

typing rules

$$\Gamma \vdash \text{true} : \text{Bool} \quad (\text{T-True})$$
$$\Gamma \vdash \text{false} : \text{Bool} \quad (\text{T-False})$$
$$\frac{x : T \in \Gamma}{\Gamma \vdash x : T} \quad (\text{T-Var})$$
$$\frac{\Gamma, x : T_1 \vdash t : T_2}{\Gamma \vdash \lambda x.t : T_1 \rightarrow T_2} \quad (\text{T-Abs})$$
$$\frac{\Gamma \vdash t_1 : T_1 \rightarrow T_2 \quad \Gamma \vdash t_2 : T_1}{\Gamma \vdash t_1 t_2 : T_2} \quad (\text{T-App})$$
$$\frac{\Gamma \vdash t_1 : \text{Bool} \quad \Gamma \vdash t_2 : T \quad \Gamma \vdash t_3 : T}{\Gamma \vdash \text{if } t_1 \text{ then } t_2 \text{ else } t_3 : T} \quad (\text{T-If})$$

Hindley-Milner type system

term and values

$$t ::= x \mid t \ t' \mid \lambda x.t \mid \text{let } x = t \text{ in } t$$

types

$$T ::= \alpha \mid T \rightarrow T$$

$$S ::= T \mid \forall \vec{\alpha}.T$$

typing rules

$$\frac{x : S \in \Gamma}{\Gamma \vdash x : S} \text{ (T-Var)}$$
$$\frac{\Gamma, x : T_1 \vdash t : T_2}{\Gamma \vdash \lambda x.t : T_1 \rightarrow T_2} \text{ (T-Abs)}$$
$$\frac{\Gamma \vdash t_1 : T_1 \rightarrow T_2 \quad \Gamma \vdash t_2 : T_1}{\Gamma \vdash t_1 \ t_2 : T_2} \text{ (T-App)}$$
$$\frac{\Gamma \vdash t_1 : S \quad \Gamma, x : S \vdash t_2 : T}{\Gamma \vdash \text{let } x = t_1 \text{ in } t_2 : T} \text{ (T-Let)}$$
$$\frac{\Gamma \vdash t : S' \quad S' \sqsubseteq S}{\Gamma \vdash t : S} \text{ (T-Inst)}$$
$$\frac{\Gamma \vdash t : S \quad \alpha \notin FV(\Gamma)}{\Gamma \vdash t : \forall \alpha.S} \text{ (T-Gen)}$$

System F

term and values

$$t ::= x \mid t \ t' \mid \lambda x : T.t \mid \Lambda \alpha.t \mid t [T]$$

$$v ::= \lambda x : T.t \mid \alpha.t$$

types

$$T ::= \alpha \mid T \rightarrow T \mid \forall \vec{\alpha}.T$$

typing rules (*in addition to* $\lambda \rightarrow$)

$$\frac{\Gamma \vdash t : T \quad \alpha \notin FV(\Gamma)}{\Gamma \vdash \Lambda \alpha.t : \forall \alpha.T} \text{ (T-TAbs)} \quad \text{(T-Gen)}$$
$$\frac{\Gamma \vdash t : \forall \alpha.T_1}{\Gamma \vdash t [T_2] : \{T_2/\alpha\}T_1} \text{ (T-TApp)} \quad \text{(T-Inst)}$$

evaluation rules (*in addition to* $\lambda \rightarrow$)

$$\frac{t \rightarrow t'}{t [T] \rightarrow t' [T]} \text{ (E-Tapp)}$$
$$(\Lambda \alpha.t) [T] \rightarrow \{T/\alpha\}t \text{ (E-TappTabs)}$$