

## Convex hull in 2D




## Before we start ...

- Link to study materials of Geometric algorithms course
- https://is.muni.cz/auth/do/sci/UMS/el/geome tricke-alg/index.html


## Convex hull

- Set $M$ is convex if a line segment connecting its two arbitrary points fully lies inside $M$

- Convex hull of a set of points $X$ in the Euclidean space corresponds to the smallest convex set containing $X$


## Convex hull

- Input: $n$ points on a plane



## Convex hull

- Convex hull in 2D:
= convex polygon
- Represented by an ordered sequence of vertices (counter clockwise)
- Convex hull in 3D:
= convex polyhedron
- Represented by a planar graph


## Convex hull - algorithms

- Gift Wrapping (Jarvis March)
- Graham Scan
- Incremental algorithm
- Divide and conquer


## Gift wrapping (Jarvis March)

- Resembles wrapping gifts, proposed by Jarvis (1973)
- Simple implementation and extension to 3D
- Assumption: set $X$ does not contain three colinear points
- Complexity: Preprocessing $O(n)$, algorithm $O\left(n^{2}\right)$


## Gift wrapping (Jarvis March)

- Principle:
- Find pivot $q\left(q=\max \left(y_{j}\right)\right)$
- Add q to the convex hull $H$
$-p_{j-1}=$ arbitrary point on $\times$ axis, $p_{j}=q, p_{i}=p_{j-1}$
- Repeat until $p_{i} \neq q$ :
- Repeat $\forall p_{i} \notin H$ and points $p_{j-1}, p_{j}$ :
- Find $p_{i}$ for that the angle $\Theta=\min \left(\Theta_{i}\right)$
- Add $p_{i}$ to $H$
- $p_{j-1}=p_{j}, p_{j}=p_{i}$


## Gift wrapping (Jarvis March)



## Gift wrapping (Jarvis March)



## Implementation

- We find point $P$ with the highest $x$-axis value this is one of the vertices of the convex hull
- In this point $P$ we determine so called separating line (often parallel to $y$ axis). All points in the input set lie in the same halfplane, determined by the separating line



## Implementation

- From P we shoot rays heading to all other points of the input set



## Implementation

- We select a ray which has the minimal angle with the first (separating) line. We have next vertext of the convex hull (2)



## Implementation

- New edge of the convex hull is 1-2



## Implementation

- Repeat this until we will reach the first point $P$ again



## Implementation





