

Distributed Computing with MapReduce

Lecture 2 of NoSQL Databases (PA195)

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Agenda



- Distributed Data Processing
- Google MapReduce
 - Motivation and History
 - Google File System (GFS)
 - MapReduce: Schema, Example, MapReduce Framework
- Apache Hadoop
 - Hadoop Modules and Related Projects
 - Hadoop Distributed File System (HDFS)
 - Hadoop MapReduce
- MapReduce in Other Systems

Distributed Data Processing



What is the **best way** of doing **distributed** processing?

Centralized (and in memory)

Don't do it, if don't have to

Big Data Processing



- Big Data analytics (or data mining)
 - need to process large data volumes quickly
 - want to use computing cluster instead of a super-computer
- Communication (sending data) between compute nodes is expensive
- => model of "moving the computing to data"

Big Data Processing II



Computing cluster architecture:



racks with compute nodes

• HW failures are rather a rule than an exception:

- 1. Files must be stored redundantly
 - over different racks to overcome also rack failures
- 2. Computations must be divided into independent tasks
 - that can be restarted in case of a failure

source: J. Leskovec, A. Rajaraman, and J. D. Ullman, Mining of Massive Datasets. 2014.

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PageRank



PageRank works by counting the number and quality of links to a page to determine a rough estimate of how important the website is.

The underlying assumption is that more important websites are likely to receive more links from other websites.



MapReduce: Origins



- In 2003, Google had the following problem:
 - How to rank tens of billions of webpages by their "importance" (PageRank) in a "reasonable" amount of time?
 - 2. How to compute these rankings efficiently when the data is scattered across thousands of computers?
- Additional factors:
 - 1. Individual data files can be enormous (terabyte or more)
 - 2. The files were rarely updated
 - the computations were read-heavy, but not very write-heavy
 - If writes occurred, they were appended at the end of the file

Google Solution



- Google found the following solutions:
 - Google File System (GFS)
 - A distributed file system
 - o MapReduce
 - A programming model for distributed data processing

Google File System (GFS)



- Files are divided into chunks (typically 64 MB)
 - The chunks are replicated at three different machines
 - ...in an "intelligent" fashion, e.g. never all on the same computer rack
 - The chunk size and replication factor are tunable
- One machine is a master, the other chunkservers
 - The master keeps track of all file metadata
 - mappings from files to chunks and locations of the chunks
 - To find a file chunk, client queries the master, and then contacts the relevant chunkservers
 - The master's metadata files are also replicated



GFS: Schema





Figure 1: GFS Architecture

MapReduce (1)



- MapReduce is a programming model sitting on the top of a Distributed File System
 Originally: no data model – data stored directly in files
- A distributed computational task has three phases:
 - 1. The map phase: data transformation
 - 2. The grouping phase
 - done automatically by the MapReduce Framework
 - 3. The reduce phase: data aggregation
- User must define only map & reduce functions





- Map function simplifies the problem in this way:
 - Input: a single data item (e.g., line of text) from a data file
 - O Output: zero or more (key, value) pairs
- The keys are not typical "keys":
 - They do not have to be unique
 - A map task can produce several key-value pairs with the same key (even from a single input)
- Map phase applies the map function to all items.





input data

map function

output data (color indicates key)

Grouping Phase



- Grouping (Shuffling): The key-value outputs from the map phase are grouped by key
 - Values sharing the same key are sent to the same reducer.
 - These values are consolidated into a single list (key, list).
 - This is convenient for the reduce function
 - This phase is realized by the MapReduce framework.



intermediate output (color indicates key)

shuffle (grouping) phase

Reduce Phase



- Reduce: combine the values for each key
 - to achieve the final result(s) of the computational task
 - o Input: (key, value-list)
 - the value-list contains all values generated for a given key in the Map phase
 - Output: (key, value-list)
 - zero or more output records





Example: Word Count



Task: Calculate word frequency in a set of documents

```
map(String key, Text value):
    // key: document name (ignored)
    // value: content of document (words)
    foreach word w in value:
        emitIntermediate(w, 1);
```

```
reduce(String key, Iterator values):
    // key: a word
    // values: a list of counts
    int result = 0;
    foreach v in values:
        result += v;
emit(key, result);
```

Example: Word Count (2)





MapReduce: Combiner



- If the reduce function is commutative & associative
 - The values can be combined in any order and combined per partes (grouped)
 - with the same result (e.g., Word Counts)
- ...then we can do "partial reductions"
 - Apply the same reduce function right after the map phase, before shuffling and redistribution to reducer nodes
- This (optional) step is known as the combiner
 Note: it is still necessary to run the reduce phase.

Example: Word Count, Combiner



Task: Calculate word frequency in a set of documents

```
combine(String key, Iterator values):
    // key: a word
    // values: a list of local counts
    int result = 0;
    foreach v in values:
        result += v;
emit(key, result);
```





MapReduce Framework



- MapReduce framework takes care about
 - Distribution and parallelizing of the computation
 - Monitoring of the whole distributed task
 - The grouping (shuffling) phase
 - putting together intermediate results
 - **Recovering** from any failures
- User must define only map & reduce functions
 o but can define also other additional functions (see below)

MapReduce Framework (2)





source: Dean, J. & Ghemawat, S. (2004). MapReduce: Simplified Data Processing on Large Clusters

MapReduce Framework: Details



- 1. Input reader (function)
 - o defines how to read data from underlying storage
- 2. Map (phase)
 - master node prepares *M* data splits and *M* idle Map tasks
 - o pass individual splits to the Map tasks that run on workers
 - these map tasks are then running
 - when a task is finished, its intermediate results are stored
- 3. Combiner (function, optional)

• combine local intermediate output from the Map phase

MapReduce Framework: Details (2)

- 4. Partition (function)
 - o to partition intermediate results for individual Reducers
- 5. Comparator (function)
 - o sort and group the input for each Reducer
- 6. Reduce (phase)
 - master node creates *R* idle Reduce tasks on workers
 - Partition function defines a data batch for each reducer
 - each Reduce task uses Comparator to create key-values pairs
 - o function Reduce is applied on each key-values pair
- 7. Output writer (function)
 - o defines how the output key-value pairs are written out

MapReduce: Example II



Task: Calculate a graph of web links

• what pages reference () each page (backlinks)

```
map(String url, Text html):
    // url: web page URL
    // html: HTML text of the page (linearized HTML tags)
foreach tag t in html:
    if t is <a> then:
        emitIntermediate(t.href, url);
```

```
reduce(String key, Iterator values):
    // key: target URLs
    // values: a list of source URLs
emit(key, values);
```

Example II: Result



Intermediate output after Map phase:

```
("http://cnn.com", "http://cnn.com")
("http://cnn.com", "http://ihned.cz")
("http://cnn.com", "http://idnes.cz")
("http://ihned.cz", "http://idnes.cz")
("http://idnes.cz", "http://idnes.cz")
```

Intermediate result after shuffle phase (the same as output after Reduce phase):

```
("http://cnn.com", ["http://cnn.com", "http://ihned.cz", "http://idnes.cz"] )
("http://ihned.cz", [ "http://idnes.cz" ])
("http://idnes.cz", [ "http://idnes.cz" ])
```

MapReduce: Example III



Task: What are the lengths of words in the input text

• output = how many words are in the text for each length

map(String key, Text value):
 // key: document name (ignored)
 // value: content of document (words)
 foreach word w in value:
 emitIntermediate(length(w), 1);

```
reduce(Integer key, Iterator values):
    // key: a length
    // values: a list of counts
    int result = 0;
    foreach v in values:
        result += v;
emit(key, result);
```

MapReduce: Features



- MapReduce uses a "shared nothing" architecture
 Nodes operate independently, sharing no memory/disk
 Common feature of many NoSQL systems
- Data partitioned and replicated over many nodes
 - Pro: Large number of read/write operations per second
 - Con: Coordination problem which nodes have my data, and when?

Applicability of MapReduce



- MR is applicable if the problem is parallelizable
- Two problems:
 - The programming model is limited (only two phases with a given schema)
 - 2. There is no data model it works only on "data chunks"
- Google's answer to the 2nd problem was **BigTable**
 - The first column-family system (2005)
 - Subsequent systems: HBase (over Hadoop), Cassandra,...

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Apache Hadoop



- Open-source software framework
 O Implemented in Java
- Able to run applications on large clusters of commodity hardware
 - Multi-terabyte data-sets
 - o Thousands of nodes
- Derived from the idea of Google's MapReduce and Google File System

Hadoop: Modules



- Hadoop Common
 - Common support functions for other Hadoop modules
- Hadoop Distributed File System (HDFS)
 - Distributed file system
 - High-throughput access to application data
- Hadoop YARN
 - Job scheduling and cluster resource management
- Hadoop MapReduce
 - YARN-based system for parallel data processing



HDFS (Hadoop Distributed File System)



- Free and open source
- Cross-platform (pure Java)
 O Bindings for non-Java programming languages
- Highly scalable
- Fault-tolerant
 - Idea: "failure is the norm rather than exception"
 - A HDFS instance may consist of thousands of machines and each can fail
 - Detection of faults
 - o Quick, automatic recovery
- Not the best in efficiency

HDFS: Data Characteristics



- Assumes:
 - o Streaming data access
 - reading the files from the beginning till the end
 - Batch processing rather than interactive user access
- Large data sets and files
- Write-once / read-many
 - A file once created does not need to be changed often
 - This assumption simplifies coherency
- Optimal applications for this model: MapReduce, web-crawlers, data warehouses, ...
HDFS: Basic Components



- Master/slave architecture
- HDFS exposes file system namespace
 File is internally split into blocks
- NameNode master server
 - Manages the file system namespace
 - Opening/closing/renaming files and directories
 - Regulates file accesses
 - Determines mapping of blocks to DataNodes
- DataNode manages file blocks
 - Block read/write/creation/deletion/replication
 - o Usually one per physical node

HDFS: Schema

HDFS Architecture





HDFS: NameNode



- NameNode has a structure called FsImage
 - Entire file system namespace + mapping of blocks to files
 + file system properties
 - Stored in a file in NameNode's local file system
 - Designed to be compact
 - Loaded in NameNode's memory (4 GB of RAM is sufficient)
- NameNode uses a transaction log called EditLog
 - to record every change to the file system's meta data
 - E.g., creating a new file, change in replication factor of a file, ..
 - EditLog is stored in the NameNode's local file system

HDFS: DataNode



- Stores data in files on its local file system
 - Each HDFS block in a separate file
 - Has no knowledge about HDFS file system
- When the DataNode starts up:
 - It generates a list of all HDFS blocks = BlockReport
 - o It sends the report to NameNode

HDFS: Blocks & Replication



- HDFS can store very large files across a cluster
 - Each file is a sequence of blocks
 - All blocks in the file are of the same size
 - Except the last one
 - Block size is configurable per file (default 128MB)
 - Blocks are replicated for fault tolerance
 - Number of replicas is configurable per file
- NameNode receives HeartBeat and BlockReport from each DataNode
 - BlockReport: list of all blocks on a DataNode

HDFS: Block Replication



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Block Replication

Namenode (Filename, numReplicas, block-ids, ...) /users/sameerp/data/part-0, r:2, {1,3}, ... /users/sameerp/data/part-1, r:3, {2,4,5}, ...



Datanodes

HDFS: Reliability



- Primary objective: to store data reliably in case of:
 - NameNode failure
 - DataNode failure
 - Network partition
 - a subset of DataNodes can lose connectivity with NameNode
- In case of absence of a HeartBeat message
 - NameNode marks "dead" the DataNodes without
 HeartBeat, and does not send any I/O requests to them.
 - The death of a DataNode typically results in re-replication

Hadoop: MapReduce



- Hadoop MapReduce requires:
 - Distributed file system (typically HDFS)
 - Engine that can distribute, coordinate, monitor and gather the results (typically YARN)
- Two main components:
 - o JobTracker (master) = scheduler
 - tracks the whole MapReduce job
 - communicates with HDFS NameNode to run the task close to the data
 - TaskTracker (slave on each node) is assigned a Map or a Reduce task (or other operations)
 - Each task runs in its own JVM

Hadoop HDFS + MapReduce





source: http://bigdata.black/architecture/hadoop/what-is-hadoop/





public class Map

extends Mapper<LongWritable, Text, Text, IntWritable> {

```
private final static IntWritable one = new IntWritable(1);
private final Text word = new Text();
```

Hadoop MR: WordCount Example

public class Reduce
 extends Reducer<Text, IntWritable, Text, IntWritable> {

```
@Override
public void reduce (Text key, Iterable<IntWritable> values,
    Context context) throws ... {
    int sum = 0;
    for (IntWritable val : values) {
        sum += val.get();
    }
    context.write(key, new IntWritable(sum));
}
```



Apache Hadoop Ecosystem

Ambari Provisioning, Managing and Monitoring Hadoop Clusters Machine Learning Data Exchange R Connectors Sqoop Mahout Workflow SQL Query Statistics Scripting Oozie Hive Columnar Store Pig Hbase YARN Map Reduce v2 Zookeeper Coordination **Distributed Processing Framework** Log Collector Flume HDFS Hadoop Distributed File System

Hadoop: Related Projects



- Avro: a data serialization system
- HBase: scalable distributed column-family database
- Cassandra: scalable distributed column-family database
- ZooKeeper: high-performance coordination service for distributed applications
- Hive: data warehouse: ad hoc querying & data summarization
- Mahout: scalable machine learning and data mining library

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MapReduce: Implementation











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Apache Spark



- Engine for distributed data processing
 - Runs over Hadoop Yarn, Apache Mesos, standalone, ...
 - Can access data from HDFS, Cassandra, HBase, AWS S3
- Can do MapReduce
 - Is much faster than pure Hadoop
 - They say 10x on the disk, 100x in memory
 - The main reason: intermediate data in memory
- Different languages to write MapReduce tasks
 o Java, Scala, Python, R

homepage: http://spark.apache.org/

Apache Spark: Example



- Example of a MapReduce task in Spark Shell
 - The shell works with Scala language
 - Example: Word count

```
val textFile = sc.textFile("hdfs://...")
val counts = textFile.flatMap(line => line.split(" "))
                .map(word => (word, 1))
                .reduceByKey(_+_)
counts.saveAsTextFile("hdfs://...")
```

• Comparison of Hadoop and Spark: <u>link</u>

MapReduce in MongoDB



```
collection "accesses":
{
    "user_id": <ObjectId>,
    "login_time": <time_the_user_entered_the_system>,
    "logout_time": <time_the_user_left_the_system>,
    "access_type": <type_of_the_access>
}
```

How much time did each user spend logged in
 Counting just accesses of type "regular"

```
db.accesses.mapReduce(
  function() { emit (this.user_id, this.logout_time - this.login_time); },
  function(key, values) { return Array.sum( values ); },
  {
    query: { access_type: "regular" },
    out: "access_times"
  }
```

References



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