

Chapter 6: Database Design Using the E-R Model

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Design Phases

- Initial phase -- characterize fully the data needs of the prospective database users.
- Second phase -- choosing a data model
 - Applying the concepts of the chosen data model
 - Translating these requirements into a conceptual schema of the database.
 - A fully developed conceptual schema indicates the functional requirements of the enterprise.
 - Describe the kinds of operations (or transactions) that will be performed on the data.



Design Phases (Cont.)

- Final Phase -- Moving from an abstract data model to the implementation of the database
 - Logical Design Deciding on the database schema.
 - Database design requires that we find a "good" collection of relation schemas.
 - Business decision What attributes should we record in the database?
 - Computer Science decision What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
 - Physical Design Deciding on the physical layout of the database



Design Alternatives

- In designing a database schema, we must ensure that we avoid two major pitfalls:
 - Redundancy: a bad design may result in repeat information.
 - Redundant representation of information may lead to data inconsistency among the various copies of information
 - Incompleteness: a bad design may make certain aspects of the enterprise difficult or impossible to model.
- Avoiding bad designs is not enough. There may be a large number of good designs from which we must choose.



Design Approaches

- Entity Relationship Model (covered in this chapter)
 - Models an enterprise as a collection of *entities* and *relationships*
 - Entity: a "thing" or "object" in the enterprise that is distinguishable from other objects
 - Described by a set of *attributes*
 - Relationship: an association among several entities
 - Represented diagrammatically by an *entity-relationship diagram*:
- Normalization Theory (Next chapter)
 - Formalize what designs are bad, and test for them



Outline of the ER Model



Entity Sets

- An entity is an object that exists and is distinguishable from other objects.
 - Example: specific person, company, event, plant
- An entity set is a set of entities of the same type that share the same properties.
 - Example: set of all persons, companies, trees, holidays
- An entity is represented by a set of attributes; i.e., descriptive properties possessed by all members of an entity set.
 - Example:

```
instructor = (ID, name, salary)
course= (course_id, title, credits)
```

A subset of the attributes form a primary key of the entity set; i.e., uniquely identifying each member of the set.



Representing Entity sets in ER Diagram

- Entity sets can be represented graphically as follows:
 - Rectangles represent entity sets.
 - Attributes listed inside entity rectangle
 - Underline indicates primary key attributes

instructor	
<u>ID</u>	
name	
salary	





Relationship Sets

• A **relationship** is an association among several entities

Example:

44553 (Peltier)	<u>advisor</u>	22222 (<u>Einstein)</u>
student entity	relationship set	instructor entity

A relationship set is a mathematical relation among n ≥ 2 entities, each taken from entity sets

 $\{(e_1, e_2, \dots, e_n) \mid e_1 \in E_1, e_2 \in E_2, \dots, e_n \in E_n\}$

where $(e_1, e_2, ..., e_n)$ is a relationship

• Example:

 $(44553,22222) \in advisor$



Relationship Sets (Cont.)

- Example: we define the relationship set *advisor* to denote the associations between students and the instructors who act as their advisors.
- Pictorially, we draw a line between related entities.





Representing Relationship Sets via ER Diagrams

Diamonds represent relationship sets.





Relationship Sets (Cont.)

- An attribute can also be associated with a relationship set.
- For instance, the *advisor* relationship set between entity sets *instructor* and *student* may have the attribute *date* which tracks when the student started being associated with the advisor





Relationship Sets with Attributes





Roles

- Entity sets of a relationship need not be distinct
 - Each occurrence of an entity set plays a "role" in the relationship
- The labels "course_id" and "prereq_id" are called **roles**.





Degree of a Relationship Set

- Binary relationship
 - involve two entity sets (or degree two).
 - most relationship sets in a database system are binary.
- Relationships between more than two entity sets are rare. Most relationships are binary. (More on this later.)
 - Example: *students* work on research *projects* under the guidance of an *instructor*.
 - relationship proj_guide is a ternary relationship between instructor, student, and project



Non-binary Relationship Sets

- Most relationship sets are binary
- There are occasions when it is more convenient to represent relationships as non-binary.
- E-R Diagram with a Ternary Relationship





Complex Attributes

- Attribute types:
 - Simple and composite attributes.
 - Single-valued and multivalued attributes
 - Example: multivalued attribute: phone_numbers
 - **Derived** attributes
 - Can be computed from other attributes
 - Example: age, given date_of_birth
- Domain the set of permitted values for each attribute



Composite Attributes

Composite attributes allow us to divide attributes into subparts (other attributes).





Representing Complex Attributes in ER Diagram

instructor
<u>ID</u>
name
first_name
middle_initial
last_name
address
street
street_number
street_name
apt_number
city
state
zip
{ phone_number }
date_of_birth
age ()



Mapping Cardinality Constraints

- Express the number of entities to which another entity can be associated via a relationship set.
- Most useful in describing binary relationship sets.
- For a binary relationship set the mapping cardinality must be one of the following types:
 - One to one
 - One to many
 - Many to one
 - Many to many



Mapping Cardinalities







One to many

Note: Some elements in *A* and *B* may not be mapped to any elements in the other set



Mapping Cardinalities







(b)

В

 b_1

 b_2

 b_3

 b_4

Note: Some elements in A and B may not be mapped to any elements in the other set



Representing Cardinality Constraints in ER Diagram

- We express cardinality constraints by drawing either a directed line (→), signifying "one," or an undirected line (—), signifying "many," between the relationship set and the entity set.
- One-to-one relationship between an *instructor* and a *student* :
 - A student is associated with at most one *instructor* via the relationship *advisor*
 - A *student* is associated with at most one *department* via *stud_dept*





One-to-Many Relationship

- one-to-many relationship between an *instructor* and a *student*
 - an instructor is associated with several (including 0) students via advisor
 - a student is associated with at most one instructor via advisor,





Many-to-One Relationships

- In a many-to-one relationship between an *instructor* and a *student*,
 - an instructor is associated with at most one student via *advisor*,
 - and a student is associated with several (including 0) instructors via advisor





Many-to-Many Relationship

- An instructor is associated with several (possibly 0) students via *advisor*
- A student is associated with several (possibly 0) instructors via advisor





Total and Partial Participation

 Total participation (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set



participation of student in advisor relation is total

- every student must have an associated instructor
- Partial participation: some entities may not participate in any relationship in the relationship set
 - Example: participation of *instructor* in *advisor* is partial



Notation for Expressing More Complex Constraints

- A line may have an associated minimum and maximum cardinality, shown in the form *I...h*, where *I* is the minimum and *h* the maximum cardinality
 - A minimum value of 1 indicates total participation.
 - A maximum value of 1 indicates that the entity participates in at most one relationship
 - A maximum value of * indicates no limit.
- Example



 Instructor can advise 0 or more students. A student must have 1 advisor; cannot have multiple advisors



Cardinality Constraints on Ternary Relationship

- We allow at most one arrow out of a ternary (or greater degree) relationship to indicate a cardinality constraint
- For example, an arrow from proj_guide to instructor indicates each student has at most one guide for a project
- If there is more than one arrow, there are two ways of defining the meaning.
 - For example, a ternary relationship *R* between *A*, *B* and *C* with arrows to *B* and *C* could mean
 - 1. Each A entity is associated with a unique entity from B and C or
 - 2. Each pair of entities from (A, B) is associated with a unique *C* entity, and each pair (A, C) is associated with a unique *B*
 - Each alternative has been used in different formalisms
 - To avoid confusion we outlaw more than one arrow



Primary Key

- Primary keys provide a way to specify how entities and relations are distinguished. We will consider:
 - Entity sets
 - Relationship sets.
 - Weak entity sets



Primary key for Entity Sets

- By definition, individual entities are distinct.
- From a database perspective, the differences among them must be expressed in terms of their attributes.
- The values of attribute values of an entity must be such that they can uniquely identify the entity.
 - No two entities in an entity set are allowed to have exactly the same value for all attributes.
- A key for an entity is a set of attributes that suffice to distinguish entities from each other



Primary Key for Relationship Sets

- To distinguish among various relationships in a relationship set we use the individual primary keys of the entities in the relationship set.
 - Let *R* be a relationship set involving entity sets E1, E2, .. En
 - The primary key for R consists of the union of the primary keys of entity sets E1, E2, ...En
 - If the relationship set R has attributes a1, a2, ..., am associated with it, then the primary key of R also includes the attributes a1, a2, ..., am
- Example: relationship set "advisor".
 - The primary key consists of *instructor.ID* and *student.ID*
- The choice of the primary key for a relationship set depends on the mapping cardinality of the relationship set.



Choice of Primary Key for Binary Relationship

- Many-to-Many relationships. The preceding union of the primary keys is a minimal superkey and is chosen as the primary key.
- One-to-Many relationships. The primary key of the "Many" side is a minimal superkey and is used as the primary key.
- Many-to-one relationships. The primary key of the "Many" side is a minimal superkey and is used as the primary key.
- One-to-one relationships. The primary key of either one of the participating entity sets forms a minimal superkey, and either one can be chosen as the primary key.



Expressing Weak Entity Sets

- In E-R diagrams, a weak entity set is depicted via a double rectangle.
- We underline the discriminator of a weak entity set with a dashed line.
- The relationship set connecting the weak entity set to the identifying strong entity set is depicted by a double diamond.
- Primary key for section (course_id, sec_id, semester, year)





Redundant Attributes

- Suppose we have entity sets:
 - *student*, with attributes: *ID*, *name*, *tot_cred*, *dept_name*
 - *department,* with attributes: *dept_name, building, budget*
- We model the fact that each student has an associated department using a relationship set stud_dept
- The attribute dept_name in student below replicates information present in the relationship and is therefore redundant
 - and needs to be removed.
- BUT: when converting back to tables, in some cases the attribute gets reintroduced, as we will see later.



(a) Incorrect use of attribute

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E-R Diagram for a University Enterprise



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Reduction to Relation Schemas



Reduction to Relation Schemas

- Entity sets and relationship sets can be expressed uniformly as *relation* schemas that represent the contents of the database.
- A database that conforms to an E-R diagram can be represented by a collection of schemas.
- For each entity set and relationship set there is a unique schema that is assigned the name of the corresponding entity set or relationship set.
- Each schema has a number of columns (generally corresponding to attributes), which have unique names.



Representing Entity Sets

• A strong entity set reduces to a schema with the same attributes

student(ID, name, tot_cred)

 A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set

section (<u>course_id, sec_id, sem, year</u>)

Example





Representation of Entity Sets with Composite Attributes



- Composite attributes are flattened out by creating a separate attribute for each component attribute
 - Example: given entity set *instructor* with composite attribute *name* with component attributes *first_name* and *last_name* the schema corresponding to the entity set has two attributes *name_first_name* and *name_last_name*
 - Prefix omitted if there is no ambiguity (name_first_name could be first_name)
- Ignoring multivalued attributes, extended instructor schema is
 - instructor(ID, first_name, middle_initial, last_name, street_number, street_name, apt_number, city, state, zip_code, date_of_birth)



Representation of Entity Sets with Multivalued Attributes

- A multivalued attribute *M* of an entity *E* is represented by a separate schema *EM*
- Schema *EM* has attributes corresponding to the primary key of *E* and an attribute corresponding to the multivalued attribute *M*
- Example: Multivalued attribute phone_number of instructor is represented by a schema: inst_phone= (<u>ID, phone_number</u>)
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema EM
 - For example, an *instructor* entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples: (22222, 456-7890) and (22222, 123-4567)



Representing Relationship Sets

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set advisor

 $advisor = (\underline{s \ id, i \ id})$





Redundancy of Schemas

- Many-to-one and one-to-many relationship sets that are total on the manyside can be represented by adding an extra attribute to the "many" side, containing the primary key of the "one" side
- Example: Instead of creating a schema for relationship set *inst_dept*, add an attribute *dept_name* to the schema arising from entity set *instructor*
- Example





Redundancy of Schemas (Cont.)

- For one-to-one relationship sets, either side can be chosen to act as the "many" side
 - That is, an extra attribute can be added to either of the tables corresponding to the two entity sets
- If participation is *partial* on the "many" side, replacing a schema by an extra attribute in the schema corresponding to the "many" side could result in null values



Redundancy of Schemas (Cont.)

- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
- Example: The section schema already contains the attributes that would appear in the sec_course schema





Extended E-R Features



Specialization

- Top-down design process; we designate sub-groupings within an entity set that are distinctive from other entities in the set.
- These sub-groupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a *triangle* component labeled ISA (e.g., *instructor* "is a" *person*).
- Attribute inheritance a lower-level entity set inherits all the attributes and relationship participations of the higher-level entity set to which it is linked.



Specialization Example

- **Overlapping** *employee* and *student*
- Disjoint instructor and secretary
- Total and partial





Representing Specialization via Schemas

- Method 1:
 - Form a schema for the higher-level entity
 - Form a schema for each lower-level entity set, including the primary key of the higher-level entity set and local attributes

schema	attributes
person	ID, name, street, city
student	ID, tot_cred
employee	ID, salary

 Drawback: getting information about, an *employee* requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema



Representing Specialization as Schemas (Cont.)

- Method 2:
 - Form a schema for each entity set with all local and inherited attributes

schema	attributes
person	ID, name, street, city
student	ID, name, street, city, tot_cred
employee	ID, name, street, city, salary

• Drawback: *name, street,* and *city* may be stored redundantly for people who are both students and employees



Generalization

- A bottom-up design process combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.



Completeness constraint

- Completeness constraint -- specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
 - total: an entity must belong to one of the lower-level entity sets
 - partial: an entity need not belong to any of the lower-level entity sets



Aggregation

- Consider the ternary relationship *proj_guide*, which we saw earlier
- Suppose we want to record evaluations of a student by a guide on a project





Aggregation (Cont.)

- Relationship sets eval_for and proj_guide represent overlapping information
 - Every *eval_for* relationship corresponds to a *proj_guide* relationship
 - However, some proj_guide relationships may not correspond to any eval_for relationships
 - So we can't discard the proj_guide relationship
- Eliminate this redundancy via *aggregation*
 - Treat relationship as an abstract entity
 - Allows relationships between relationships
 - Abstraction of relationship into new entity



Aggregation (Cont.)

- Eliminate this redundancy via aggregation without introducing redundancy, the following diagram represents:
 - A student is guided by a particular instructor on a particular project
 - A student, instructor, project combination may have an associated evaluation





Reduction to Relational Schemas

- To represent aggregation, create a schema containing
 - Primary key of the aggregated relationship,
 - The primary key of the associated entity set
 - Any descriptive attributes
- In our example:
 - The schema *eval_for* is:

eval_for (s_ID, project_id, i_ID, evaluation_id)

• The schema *proj_guide* is redundant.



Design Issues



Entities vs. Attributes

Use of entity sets vs. attributes



 Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)



Entities vs. Relationship sets

Use of entity sets vs. relationship sets

A possible guideline is to designate a relationship set to describe an action that occurs between entities



Placement of relationship attributes

For example, attribute date as attribute of advisor or as attribute of student



Binary Vs. Non-Binary Relationships

- Although it is possible to replace any non-binary (*n*-ary, for *n* > 2) relationship set by a number of distinct binary relationship sets, a *n*-ary relationship set shows more clearly that several entities participate in a single relationship.
- Some relationships that appear to be non-binary may be better represented using binary relationships
 - For example, a ternary relationship *parents*, relating a child to his/her father and mother, is best replaced by two binary relationships, *father* and *mother*
 - Using two binary relationships allows partial information (e.g., only the mother being known)
 - But there are some relationships that are naturally non-binary
 - Example: proj_guide



Converting Non-Binary Relationships to Binary Form

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
 - Replace *R* between entity sets A, B and C by an entity set *E*, and three relationship sets:
 - 1. R_A , relating E and A3. R_C , relating E and C2. R_B , relating E and B
 - Create an identifying attribute for *E* and add any attributes of *R* to *E*
 - For each relationship (*a_i*, *b_i*, *c_i*) in *R*, create
 - 1. a new entity e_i in the entity set E 2. add (e_i, a_i) to R_A
 - 3. add (e_i, b_i) to R_B





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E-R Design Decisions

- The use of an attribute or entity set to represent an object.
- Whether a real-world concept is best expressed by an entity set or a relationship set.
- The use of a ternary relationship versus a pair of binary relationships.
- The use of a strong or weak entity set.
- The use of specialization/generalization contributes to modularity in the design.
- The use of aggregation can treat the aggregate entity set as a single unit without concern for the details of its internal structure.



Summary of Symbols Used in E-R Notation



identifying relationship set for weak entity set



attributes: simple (A1), composite (A2) and multivalued (A3) derived (A4)









total participation of entity set in relationship

E	
A1	

discriminating attribute of weak entity set



Symbols Used in E-R Notation (Cont.)



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End of Chapter 6