



CEITEC

Central European Institute of Technology
BRNO | CZECH REPUBLIC

**Introduction to Bioinformatics
(LF:DSIB01)**

**Week 5 : Dynamic
Programming**



Sequence Alignment – Hamming Distance (1950)

- Given two equally sized sequences, what is the distance between the sequences?

Hamming Distance = number of mismatches

ATCAGGTCATGCAG

ATCAGG**AC**ATGC**AC**

Hamming Distance fails with insertion

ATCAGGTCATGCAG

At**TCAGGAC**ATGC**A**

Sequence Alignment – Levenshtein Distance (1965)

Given two potentially related sequences (of any length), what is the distance between the sequences?

Distance = minimum number of edits needed to get from A to B

Edits = **Substitution** of letter, **addition** of letter, **deletion** of letter

A = “ACGTCATCA”

B = “TAGTGTCA”

Scoring Function

Score(AB) = Total cost of editing A to B

Cost of substitution (mutation) (-1)

Cost of indel (-1)

Reward for match (+1)

What is the best alignment (highest score)?

How do you calculate it?

Sequence Alignment – Levenshtein Distance (1965)

Score = -5

A C G T C A T C A

T A G T G T C A _

Sequence Alignment – Levenshtein Distance (1965)

Score = -5

A C G T C A T C A
T A G T G T C A _

Score = +1

A C G T C A T C A
T A G T G _ T C A

Sequence Alignment – Levenshtein Distance (1965)

Score = -5

A C **G** **T** C A T C A
T A **G** **T** G T C A _

Score = +1

A C **G** **T** C A **T** **C** **A**
T A **G** **T** G _ **T** **C** **A**

Score = +2

_ **A** C **G** **T** C A **T** **C** **A**
T **A** _ **G** **T** G _ **T** **C** **A**

Sequence Alignment

Way too many sequences for
Brute Force enumeration

$$\binom{n+m}{m} \frac{(m+n)!}{(m!)^2} \approx \frac{2^{m+n}}{\sqrt{\pi m}}$$

n	Enumeration
10	184,756
20	1.40E+11
100	9.00E+58

Coin Change (Euro)

Convert an amount of money into the fewest number of coins

Input: Amount of money (M)

Output: the smallest number of 50c (a), 20c (b), 10c (c), 5c (d), 2c (e) and 1c (f) such that $50a+20b+10c+5d+2e+1f = M$

```
1  while  $M > 0$ 
2       $c \leftarrow$  Largest coin that is smaller than (or equal to)  $M$ 
3      Give coin with denomination  $c$  to customer
4       $M \leftarrow M - c$ 
```

Try: $M=60c$; $M=55c$; $M=40c$

Coin Change (Generalised)

Input: An amount of money M , and an array of d denominations $\mathbf{c} = (c_1, c_2, \dots, c_d)$, in decreasing order of value ($c_1 > c_2 > \dots > c_d$).

Output: A list of d integers i_1, i_2, \dots, i_d such that $c_1 i_1 + c_2 i_2 + \dots + c_d i_d = M$, and $i_1 + i_2 + \dots + i_d$ is as small as possible.

Coin Change (US)

Input: An amount of money M , and an array of d denominations $\mathbf{c} = (c_1, c_2, \dots, c_d)$, in decreasing order of value ($c_1 > c_2 > \dots > c_d$).

Output: A list of d integers i_1, i_2, \dots, i_d such that $c_1 i_1 + c_2 i_2 + \dots + c_d i_d = M$, and $i_1 + i_2 + \dots + i_d$ is as small as possible.

BETTERCHANGE(M, \mathbf{c}, d)

```
1   $r \leftarrow M$ 
2  for  $k \leftarrow 1$  to  $d$ 
3       $i_k \leftarrow r/c_k$ 
4       $r \leftarrow r - c_k \cdot i_k$ 
5  return  $(i_1, i_2, \dots, i_d)$ 
```

NB: Division = "floor"

Try

$M = 40$; $c_1=25$, $c_2=10$, $c_3=5$, $c_4=1$



Coin Change (US)

Input: An amount of money M , and an array of d denominations $\mathbf{c} = (c_1, c_2, \dots, c_d)$, in decreasing order of value ($c_1 > c_2 > \dots > c_d$).

Output: A list of d integers i_1, i_2, \dots, i_d such that $c_1 i_1 + c_2 i_2 + \dots + c_d i_d = M$, and $i_1 + i_2 + \dots + i_d$ is as small as possible.

BETTERCHANGE(M, \mathbf{c}, d)

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5 return  $(i_1, i_2, \dots, i_d)$ 
```

NB: Division = "floor"



$M = 40$; $c_1=25$, $c_2=20$, $c_3=10$, $c_4=5$, $c_5=1$



Discontinued
1875
for being too
confusing

Pseudocode Exercise: Coin Change (US coins)

Input: An amount of money M , and an array of d denominations $\mathbf{c} = (c_1, c_2, \dots, c_d)$, in decreasing order of value ($c_1 > c_2 > \dots > c_d$).

Output: A list of d integers i_1, i_2, \dots, i_d such that $c_1 i_1 + c_2 i_2 + \dots + c_d i_d = M$, and $i_1 + i_2 + \dots + i_d$ is as small as possible.

BETTERCHANGE(M, \mathbf{c}, d)

```
1  $r \leftarrow M$ 
2 for  $k \leftarrow 1$  to  $d$ 
3      $i_k \leftarrow r/c_k$ 
4      $r \leftarrow r - c_k \cdot i_k$ 
5 return  $(i_1, i_2, \dots, i_d)$ 
```

BetterChange

40=
1x25 + 1x10 + 1x5=
3 coins

Incorrect!
40 = 2x20=
2 coins



$M = 40$; $c_1=25$, $c_2=20$, $c_3=10$, $c_4=5$, $c_5=1$



Discontinued
1875
for being too
confusing

NB: Division = "floor"

Pseudocode Exercise: Coin Change

Input: An amount of money M , and an array of d denominations $\mathbf{c} = (c_1, c_2, \dots, c_d)$, in decreasing order of value ($c_1 > c_2 > \dots > c_d$).

Output: A list of d integers i_1, i_2, \dots, i_d such that $c_1i_1 + c_2i_2 + \dots + c_di_d = M$, and $i_1 + i_2 + \dots + i_d$ is as small as possible.

Tries every combination
Guaranteed to find optimal
Slow

Brute Force Algorithm

Input: An amount of money M , and an array of d denominations $\mathbf{c} = (c_1, c_2, \dots, c_d)$, in decreasing order of value ($c_1 > c_2 > \dots > c_d$).

Output: A list of d integers i_1, i_2, \dots, i_d such that $c_1i_1 + c_2i_2 + \dots + c_di_d = M$, and $i_1 + i_2 + \dots + i_d$ is as small as possible.

BRUTEFORCECHANGE(M, \mathbf{c}, d)

```
1  smallestNumberOfCoins  $\leftarrow \infty$ 
2  for each  $(i_1, \dots, i_d)$  from  $(0, \dots, 0)$  to  $(M/c_1, \dots, M/c_d)$ 
3      valueOfCoins  $\leftarrow \sum_{k=1}^d i_k c_k$ 
4      if valueOfCoins =  $M$ 
5          numberOfCoins  $\leftarrow \sum_{k=1}^d i_k$ 
6          if numberOfCoins < smallestNumberOfCoins
7              smallestNumberOfCoins  $\leftarrow$  numberOfCoins
8              bestChange  $\leftarrow (i_1, i_2, \dots, i_d)$ 
9  return (bestChange)
```

Tries every combination
Guaranteed to find optimal
Slow

Dynamic Programming (Bellman, 1954)

1 cent, 5 cents, 10 cents

1	2	3	4	5	6	7	8	9	10	11
1	2	3	4	1	2	3	4	5	1	2
12	13	14	15	16	17	18	19	20	21	22
3	4	5	2	?						

Value of n

Pick the lowest number from:

$$1 + \text{Val}(n-1)$$

$$1 + \text{Val}(n-5)$$

$$1 + \text{Val}(n-10)$$

Val(16)

Min()

$$1 + \text{Val}(15) = 1 + 2 = 3$$

$$1 + \text{Val}(11) = 1 + 2 = 3$$

$$1 + \text{Val}(6) = 1 + 2 = 3$$

$$\text{Val}(16) = 3$$

$$\text{Val}(17) = ?$$

Dynamic Programming

Min

V(39)

V(35)

V(30)

V(20)

V(15)

+1

BetterChange

40=

$1 \times 25 + 1 \times 10 + 1 \times 5 =$

3 coins

Incorrect!

$40 = 2 \times 20 =$

2 coins



$M = 40; c_1=25, c_2=20, c_3=10, c_4=5, c_5=1$



Discontinued
1875
for being too
confusing

1	2	3	4	5	6	7	8	9	10
1	2	3	4	1	2	3	4	5	1
11	12	13	14	15	16	17	18	19	20
2	3	4	5	2	3	4	5	6	1
21	22	23	24	25	26	27	28	29	30
2	3	4	5	1	2	3	4	5	2
31	32	33	34	35	36	37	38	39	40
3	4	5	6	2	3	4	5	6	?

Backpropagation

What is the sequence that brought you here?

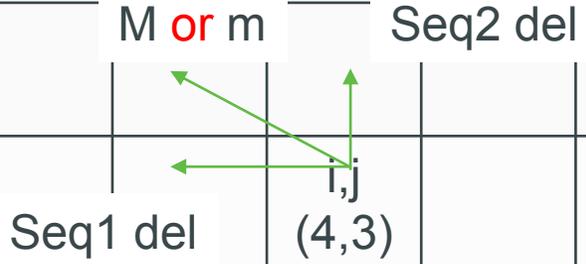
Min
 $V(39) = \text{cent}$
 $V(35) = \text{five}$
 $V(30) = \text{ten}$
 $V(20) = \text{twenty}$
 $V(15) = \text{twefive}$
 $+1$

1	2	3	4	5	6	7	8	9	10
1 (cent)	2 (cent)	3	4	1 (five)	2 (cent)	3	4	5	1 (ten)
11	12	13	14	15	16	17	18	19	20
2	3	4	5	2 (ten)	3	4	5	6	1 twenty
21	22	23	24	25	26	27	28	29	30
2	3	4	5	1 twefive	2	3	4	5	2 (ten)
31	32	33	34	35	36	37	38	39	40
3	4	5	6	2 (ten)	3	4	5	6 (cent)	?

Needleman–Wunsch algorithm (1970)

Seq1 = "ACGTCATCA"
 Seq2 = "TAGTGTCA"
 Scoring Function
 Mut / In / Del (-1)
 Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A
Seq2	0	-1	-2	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1									
A (j=2)	-2			M or m	Seq2 del					
G	-3									
T	-4									
G	-5									
T	-6									
C	-7									
A	-8									



Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A
Seq2	0	-1 (L)	-2 (L)	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1 (U)	-1 (d)	-2 (d)	-3 (d)	-2 (d)					
A (j=2)	-2 (U)									
G	-3									
T	-4									
G	-5									
T	-6									
C	-7									
A	-8									



Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A
Seq2	0	-1 (L)	-2 (L)	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1 (U)	-1 (d)	-2 (d)	-3 (d)	-2 (d)	-3	-4	-5	-6	-7
A (j=2)	-2 (U)	0 (d)	-1 (L)	-2	-3	-4	-2	-3	-4	-3
G	-3	-1	-1	0 (d)	-1	-2	-3	-3	-4	-4
T	-4	-2								
G	-5	-3								
T	-6	-4								
C	-7	-5								
A	-8	-6								



Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A
Seq2	0	-1 (L)	-2 (L)	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1 (U)	-1 (d)	-2 (d)	-3 (d)	-2 (d)	-3	-4	-5	-6	-7
A (j=2)	-2 (U)	0 (d)	-1 (L)	-2	-3	-4	-2	-3	-4	-3
G	-3	-1	-1	0 (d)	-1	-2	-3	-3	-4	-4
T	-4	-2	-2	-1	+1	0	-1	-2	-3	-4
G	-5	-3	-3	-1	0	0	-1	-2	-3	-4
T	-6	-4	-4	-2	0	-1	-1	0	-1	-2
C	-7	-5	-3	-3	-1	1	-1	-1	1	0
A	-8	-6	-4	-4	-2	0	2	1	0	2



Backpropagation

Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A (i=9)
Seq2	0	-1 (L)	-2 (L)	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1 (U)	-1 (d)	-2 (d)	-3 (d)	-2 (d)	-3	-4	-5	-6	-7
A (j=2)	-2 (U)	0 (d)	-1 (L)	-2	-3	-4	-2	-3	-4	-3
G	-3	-1	-1	0 (d)	-1	-2	-3	-3	-4	-4
T	-4	-2	-2	-1	+1	0	-1	-2	-3	-4
G	-5	-3	-3	-1	0	0	-1	-2	-3	-4
T	-6	-4	-4	-2	0	-1	-1	0	-1	-2
C	-7	-5	-3	-3	-1	1	-1	-1	1	0
A (j=8)	-8	-6	-4	-4	-2	0	2	1	0	2 (d)

[9,8] diag A A
[8,7] ?

Backpropagation

Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A (i=9)
Seq2	0	-1 (L)	-2 (L)	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1 (U)	-1 (d)	-2 (d)	-3 (d)	-2 (d)	-3	-4	-5	-6	-7
A (j=2)	-2 (U)	0 (d)	-1 (L)	-2	-3	-4	-2	-3	-4	-3
G	-3	-1	-1	0 (d)	-1	-2	-3	-3	-4	-4
T	-4	-2	-2	-1	+1	0	-1	-2	-3	-4
G	-5	-3	-3	-1	0	0	-1	-2	-3	-4
T	-6	-4	-4	-2	0	-1	-1	0	-1	-2
C	-7	-5	-3	-3	-1	1	-1	-1	1	0
A (j=8)	-8	-6	-4	-4	-2	0	2	1	0	2 (d)

[9,8] diag AA

[8,7] diag CC

[7,6] diag TT

[6,5] diag AG (m)

Backpropagation

Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

i,j	Seq1	A (i=1)	C (i=2)	G	T	C (i=5)	A	T	C	A (i=9)
Seq2	0	-1 (L)	-2 (L)	-3	-4	-5	-6	-7	-8	-9
T (j=1)	-1 (U)	-1 (d)	-2 (d)	-3 (d)	-2 (d)	-3	-4	-5	-6	-7
A (j=2)	-2 (U)	0 (d)	-1 (L)	-2	-3	-4	-2	-3	-4	-3
G	-3	-1	-1	0 (d)	-1	-2	-3	-3	-4	-4
T (4)	-4	-2	-2	-1	+1	0	-1	-2	-3	-4
G	-5	-3	-3	-1	0	0	-1	-2	-3	-4
T	-6	-4	-4	-2	0	-1	-1	0	-1	-2
C	-7	-5	-3	-3	-1	1	-1	-1	1	0
A (j=8)	-8	-6	-4	-4	-2	0	2	1	0	2 (d)

Score = 2

[9,8] diag AA
 [8,7] diag CC
 [7,6] diag TT
 [6,5] diag AG (m)
 [5,4] left C_
 [4,4] diag TT
 [3,3] diag GG
 [2,2] left C_
 [1,2] diag AA
 [0,1] up _T

 ACGTCATCA
 TA GT GTCA
 - + - + + - - + + +

Local Alignment

- Local Alignment

Score = 3

ACGTCATCA

TAGTG_TCA

Score = 4

AC__GTCA TCA

TAGTGTCA__

Smith-Waterman algorithm (1981)

Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

or ZERO

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A
Seq2	0	0	0	0	0	0	0	0	0	0
T (j=1)	0									
A (j=2)	0			M or m		Seq2 del				
G	0									
T	0									
G	0									
T	0									
C	0									
A	0									

M or m

Seq2 del



i,j
(4,3)

Seq1 del



Smith-Waterman algorithm (1981)

Seq1 = "ACGTCATCA"

Seq2 = "TAGTGTCA"

Scoring Function

Mut / In / Del (-1)

Match (+1)

or ZERO

i,j	Seq1	A (i=1)	C (i=2)	G	T	C	A	T	C	A
Seq2	0	0	0	0	0	0	0	0	0	0
T (j=1)	0	0	0	0	1	0	0	1	0	0
A (j=2)	0	1	0	0	0	0	1	0	0	1
G	0	0	0	1	0	0	0	0	0	0
T	0	0	0	0	2	0	0	1	0	0
G	0	0	0	1	1	1	0	0	0	0
T	0	0	0	0	2	1	0	1	0	0
C	0	0	1	0	1	3	2	1	2	1
A	0	0	0	0	0	2	4	3	2	3

Score: ATCAGGT vs ATTCAGG

Global Alignment

		A	T	C	A	G	G	T
A								
T								
T								
C								
A								
G								
G								

Local Alignment

		A	T	C	A	G	G	T
A								
T								
T								
C								
A								
G								
G								

Score: ATCAGGT vs ATTCAGG

Global Alignment

		A	T	C	A	G	G	T
	0	-1	-2	-3	-4	-5	-6	-7
A	-1							
T	-2							
T	-3							
C	-4							
A	-5							
G	-6							
G	-7							

Local Alignment

		A	T	C	A	G	G	T
	0	0	0	0	0	0	0	0
A	0							
T	0							
T	0							
C	0							
A	0							
G	0							
G	0							

Score: ATCAGGT vs ATTCAGG

A_TCAGGT
 ATTCAGG_
 OR
 AT_CAGGT
 ATTCAGG_

TCAGG
 TCAGG
 OR
 AT_CAGG
 ATTCAGG

Global Alignment

Local Alignment

		A	T	C	A	G	G	T
	0	-1	-2	-3	-4	-5	-6	-7
A	-1	1	0	-1	-2	-3	-4	-5
T	-2	0	2	1	0	-1	-2	-3
T	-3	-1	1	0	-1	-2	-3	-1
C	-4	-2	0	2	1	0	-1	-2
A	-5	-3	-1	1	3	2	1	0
G	-6	-4	-2	0	2	4	3	2
G	-7	-5	-3	-1	1	3	5	4

		A	T	C	A	G	G	T
	0	0	0	0	0	0	0	0
A	0	1	0	0	1	0	0	0
T	0	0	2	1	0	0	0	1
T	0	0	1	1	0	0	0	1
C	0	0	0	2	1	0	0	0
A	0	1	0	1	3	2	1	0
G	0	0	0	0	2	4	3	2
G	0	0	0	0	1	3	5	4

Binding Matrix

- Binding Matrix instead of Match/Mismatch
- How do you change the algorithm?

a

	water				
	A	C	G	U	T
ALA	0.2	0.9	1.0	0.4	0.3
ARG	-0.8	-1.1	-1.2	-0.5	0.3
ASN	-1.0	-0.7	-0.6	-0.7	-0.7
ASP	0.5	2.4	1.0	2.8	3.7
CYSH	-1.2	-0.2	-0.9	-0.8	-0.6
GLN	-2.0	-0.5	-2.2	-1.5	-1.9
GLU	2.0	1.6	-0.3	2.2	3.7
HISA	-1.8	-0.9	-1.8	-1.5	-1.9
HISB	-1.2	-0.2	-1.9	-0.6	-1.2
HISH	-0.7	0.0	-0.6	0.5	0.2
ILE	-2.8	-1.9	-2.6	-1.5	-1.6
LEU	-2.0	-1.9	-2.6	-1.8	-1.7
LYSH	0.9	1.4	-0.5	0.6	0.9
MET	-2.5	-1.4	-2.2	-2.1	-2.3
PHE	-3.3	-2.7	-3.6	-3.1	-3.9
SER	0.3	1.1	1.1	0.5	-0.4
THR	-0.5	0.6	0.2	-0.2	-0.7
TRP	-3.8	-3.6	-4.8	-4.8	-6.1
TYR	-3.9	-3.3	-5.1	-3.3	-4.3
VAL	-1.5	-1.2	-1.4	-1.0	-1.1

Extending Gap

- Creation of a gap is much more rare than extension of a gap
- Based on evolution of biological sequences
- Indel Penalty
- Indel Extension Penalty
- How do you change the implementation?



CEITEC



@CEITEC_Brno

