

EXERCISES IN CATEGORY THEORY 7

1. EQUIVALENCE OF CATEGORIES

An equivalence of categories consists of categories \mathcal{C} and \mathcal{D} together with functors $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{D} \rightarrow \mathcal{C}$ and natural isomorphisms $1_{\mathcal{C}} \rightarrow GF$ and $1_{\mathcal{D}} \rightarrow FG$.¹ We also call the functors involved $F : \mathcal{C} \rightarrow \mathcal{D}$ and $\mathcal{D} \rightarrow \mathcal{C}$ equivalences of categories.

- (1) Let Cat_1 denote the category of categories with one object and functors between them. There is a functor $U : \text{Cat}_1 \rightarrow \text{Mon}$ which sends a category \mathcal{C} with one object x to the monoid $\mathcal{C}(x, x)$. Show that U is part of an equivalence of categories.
- (2) Show that equivalence of categories is an equivalence relation on categories.

A functor $U : \mathcal{C} \rightarrow \mathcal{D}$ is said to be *essentially surjective* if given $a \in \mathcal{D}$ there exists $b \in \mathcal{C}$ such that Fb is isomorphic to a in \mathcal{D} . A functor $U : \mathcal{C} \rightarrow \mathcal{D}$ is said to be *fully faithful* if given $f : Ua \rightarrow Ub$ there exists a unique arrow $g : a \rightarrow b$ such that $Ug = f$. (In other words, the function $F_{a,b} : \mathcal{C}(a, b) \rightarrow \mathcal{D}(Fa, Fb)$ is a bijection.)

In fact a functor $U : \mathcal{C} \rightarrow \mathcal{D}$ is an equivalence of categories if and only if it is essentially surjective and fully faithful.

Use this result where useful in the following exercises.

- (3) Set_{par} is the category of sets and partial functions. A *partial function* $(U, f) : X \rightarrow Y$ consists of a subset $U \subseteq X$ and function $f : U \rightarrow Y$. Thus a function from X to Y which is only defined on U . The composite of $(U, f) : X \rightarrow Y$ and $(V, g) : Y \rightarrow Z$ is the partial function $(W, gf) : X \rightarrow Z$ where $W = \{x \in U : fx \in V\}$. (Draw some diagrams to see this).

The category of pointed sets Set_* has objects (X, x) where $x \in X$ and morphisms $f : (X, x) \rightarrow (Y, y)$ are functions such that $fx = y$.

There is a functor $F : \text{Set}_{par} \rightarrow \text{Set}_*$ defined by $FX = (X + 1, * \in 1)$ and sends $(U, f) : X \rightarrow Y$ to $k : (X + 1, *) \rightarrow (Y + 1, *)$ defined by $kx = fx$ if $x \in U$ and $k(y) = *$ otherwise. Show that F is a functor and an equivalence of categories.

- (4) Show that if $F : \mathcal{C} \rightarrow \mathcal{D}$ is an equivalence then it preserves terminal and initial objects. (In fact, an equivalence preserves any limits or colimits that exist.)
- (5) (*Harder*) Prove that a functor is an equivalence of categories if and only if it is essentially surjective on objects and fully faithful.

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¹Note that the direction of the natural isomorphisms does not matter since if $1_{\mathcal{C}} \rightarrow GF$ is a natural isomorphism then we get a natural isomorphism $GF \rightarrow 1_{\mathcal{C}}$ by taking inverses.