

Lecture 8 - There was no week 7 (suattek)

(Reminder of last time.)

- Ω a signature
- X a set, $Tr(X)$ is set defined ind. by
 - if $x \in X$ then $x \in Tr(X)$
 - if $s \in \Omega_n$ & $t_1, \dots, t_n \in Tr(X)$ then $s(t_1, \dots, t_n) \in Tr(X)$.
- $Tr(X)$ is an Ω -alg w' substitution of terms:
 - if $s \in \Omega_n, t_1, \dots, t_n \in Tr(X) \mapsto s(t_1, \dots, t_n) \in Tr(X)$
- It is the free Ω -algebra on X :

we have $\begin{matrix} X & \xrightarrow{\eta_X} & Tr(X) \\ & \searrow & \downarrow \eta \\ & & A \end{matrix}$ $\exists! \bar{v}$ an Ω -alg hom st $\bar{v} \circ \eta_X = \eta$.
 η an X -tuple of elements of A

- $\bar{v}(x) = \eta(x)$

- $\bar{v}(s(t_1, \dots, t_n)) = s^A(\bar{v}(t_1), \dots, \bar{v}(t_n))$.

- What \bar{v} does is substitutes the X -tuple η for the variables in X .

Eg. $\Omega = \{ \cdot, e \}$ for magmas

$X = \{ x, y, z \} \quad \eta = (a, b, c) \rightarrow A$

$Tr(X)$
 $\downarrow \eta$
 $(x \cdot y) \cdot z, x \cdot (y \cdot z) \rightarrow (a \cdot b) \cdot c, a \cdot (b \cdot c)$

Equations

Defⁿ) let Ω be signature, X a set. By an Ω -equation in variables X , we mean a pair $(s, t) \in \text{Tr}(X)^2$.

Remark) Often informally write an Ω -equation (s, t) as " $s = t$ ".

Example

$\Omega = (e, \cdot)$ sig. For magmas, the following are equations:

- $x \cdot e = x$
 - $e \cdot x = x$
 - $(x \cdot y) \cdot z = x \cdot (y \cdot z)$
- } in variables $\{x, y, z\}$

Defⁿ) let $(s, t) \in \text{Tr}(X)^2$ be an equation.

- An X -tuple $v: X \rightarrow A$ of elements of an Ω -algebra A satisfies the equation $s = t$

if $\bar{v}(s) = \bar{v}(t)$ where

$\bar{v}: \text{Tr}(X) \rightarrow A$ is the map described above.

- The Ω -algebra A sat. the equation $s = t$ if each X -tuple v of A sat. the equation
(We sometimes write $A \models s = t$.)

Example

$$\mathcal{R} = (\cdot, e)$$

$$(x \cdot y) \cdot z, x \cdot (y \cdot z) \in \text{Tr}(\{x, y, z\})$$

- Then $\{x, y, z\} \xrightarrow{(a, b, c)} A$ set the

$$\text{eq}^n \iff (a \cdot b) \cdot c = a \cdot (b \cdot c) \in A$$

- $A \models (x \cdot y) \cdot z = x \cdot (y \cdot z)$
iff A is associative.

Defⁿ) let \mathcal{R} be a sig., & E a set of equations. By an (\mathcal{R}, E) -algebra we mean an \mathcal{R} -algebra A such that $A \models s = t$ for each equation $(s, t) \in E$.

Examples

- For $\mathcal{R} = (\cdot, e)$ &

$$E = \left\{ \begin{array}{l} (x \cdot y) \cdot z = x \cdot (y \cdot z) \\ x \cdot e = x \\ e \cdot x = e \end{array} \right\}$$

an (\mathcal{R}, E) -alg is a monoid,

- If take

$$\mathcal{R}' = (\cdot, e, (-)^{-1}) \quad \text{unary op.}$$

$$\& E' = E \cup \{x \cdot x^{-1} = e, x^{-1} \cdot x = e\}$$

then (Ω, E) -algebra is a group.

Defⁿ) For (Ω, E) as above, we define (Ω, E) -Alg \longleftrightarrow Ω -Alg as the Full subcategory of (Ω, E) -algebras.

This means:

- obs: are (Ω, E) -algebras
- morphisms: Ω -alg. homomorphisms.
- The inclusion (Ω, E) -Alg \xrightarrow{i} Ω -Alg is then a Fully Faithful functor.
- We obtain a composite Forgetful Functor to Set, as depicted below

$$\begin{array}{ccc} (\Omega, E)\text{-Alg} & \xrightarrow{i} & \Omega\text{-Alg} \\ & \searrow \scriptstyle U & \swarrow \scriptstyle U_{\Omega} \\ & \text{Set} & \end{array}$$

- later, we will see that i & U have left adjoints.

Examples

- when I spoke of "algebraic categories" earlier in course, the precise meaning is cat. of the form

$$(\Omega, E)\text{-Alg}$$

This framework captures all

of the examples we have been talking about -

Vect, Grp, Ring, Mon, G-Set.

Goal now: study good properties of categories of the form Ω -Alg & (Ω, E) -Alg.

- Firstly (today) we look at Ω -Alg - the case of (Ω, E) -Alg follows easily from Ω -Alg.
- In particular, will study limits, kernels and quotients.

Firstly, subalgebras, homomorphic images & image factorisation.

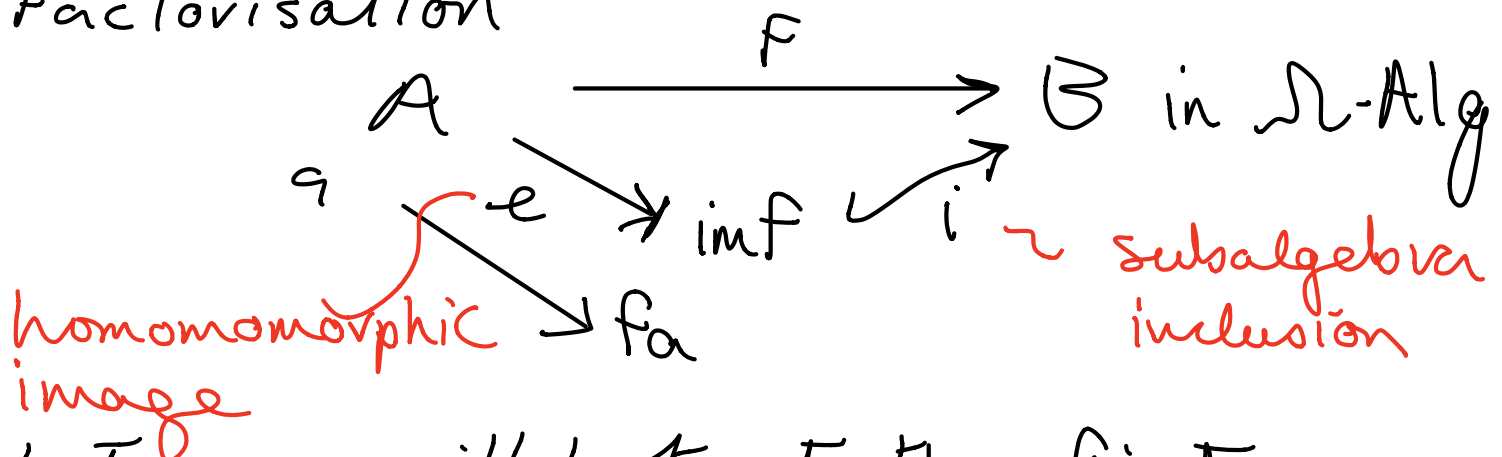
Def) - Let $A \in \Omega$ -Alg. A subalgebra $B \hookrightarrow A$ of A is a subset B of A such that - if $s \in \Omega_n$ & $b_1, \dots, b_n \in B$ then $s^A(b_1, \dots, b_n) \in B$.

- In particular, B is then an Ω -alg & the incl. $B \hookrightarrow A$ a injective homomorphism.

Defⁿ) let $A \in \Omega$ -Alg. A homomorphic image of A is a surjective homomorphism $f: A \twoheadrightarrow B$.
 means surjective \twoheadrightarrow

- Let $f: A \twoheadrightarrow B \in \Omega$ -Alg.

- Then let $\text{im}f = \{ b \in B : \exists a \in A \text{ with } fa = b \}$
- Then $\text{im}f \hookrightarrow B$ is a subalgebra of B :
indeed, if $s \in \Omega_n$, $b_1 = fa_1, \dots, b_n = fa_n$,
then $s(b_1, \dots, b_n) = s(fa_1, \dots, fa_n)$
 $= fs(a_1, \dots, a_n) \in \text{im}f$.
- In particular, we obtain a factorisation



- Later, we will look at the First isomorphism theorem which explains how to view $\text{im}f$ as a quotient.

Limits of Ω -algebras

Proposition

Ω -Alg has (infinite) products and equalisers
& $u: \Omega\text{-Alg} \longrightarrow \text{Set}$ preserves them.

(Remark: these generate all limits - not proved in course.)

Proof

- Consider a set I and Family $(A_i)_{i \in I}$

- of \mathcal{R} -algebras, (i.e. $A: I \rightarrow \mathcal{R}\text{-Alg}$)
- Their product as sets is the direct product

$$\prod_{i \in I} A_i = \{ \bar{a} = (a_i)_{i \in I} : a_i \in A_i \} \xrightarrow{p_i} A_i$$

$$(a_i)_{i \in I} \longmapsto a_i$$

- We want to show that $\prod_{i \in I} A_i$ has the structure of \mathcal{R} -algebra such that each p_i is a homomorphism:

this says given $s \in \mathcal{R}$ and $\bar{a}^1, \dots, \bar{a}^n$ we have $s(\bar{a}^1, \dots, \bar{a}^n)_i = s^i(\bar{a}^1_i, \dots, \bar{a}^n_i) \in A_i$.

In other words, we are forced to equip $\prod A_i$ with component-wise \mathcal{R} -algebra structure.

- Given \mathcal{R} -alg. B & homs $(f_i: B \rightarrow A_i)_{i \in I}$ we have a unique function

$$f: B \longrightarrow \prod_{i \in I} A_i \text{ such } p_i \circ f = f_i; \text{ that}$$

namely $(fb)_i = f_i(b)$.

- must check f is a homomorphism:

$$f s(b^1, \dots, b^n) = s(fb^1, \dots, fb^n) \in \prod_{i \in I} A_i$$

check components at $i \in I$

$$f_i s(b^1, \dots, b^n) = s((fb^1)_i, \dots, (fb^n)_i) \xrightarrow{\text{pointwise } \mathcal{R}\text{-alg. structure}} s(f_i b^1, \dots, f_i b^n)$$

$f \sim \text{hom.}$

- Given $A \xrightleftharpoons[f]{g} B$ their equaliser

is $E = \{ x \in A : Fx = gx \} \xrightarrow{i} A$ in Set.

In fact, E is a subalgebra of A :
 if $s \in \Omega_n$ & x_1, \dots, x_n st $f x_i = g x_i$ then
 $f s(x_1, \dots, x_n) = s(f x_1, \dots, f x_n) = s(g x_1, \dots, g x_n) = g s(x_1, \dots, x_n)$
f a hom assump. g hom.

In partic., $i: E \hookrightarrow A$ is a homomorphism
 and easy to check univ. prop. of the
 equaliser. \square

What about colimits?

- Key sort - quotients by congruences.

Congruences generalise: e. rels for sets

- normal subgroups of groups
- 2-sided ideals for comm rings

Def) Let A be an Ω -algebra. An equivalence
 relation $E \subseteq A^2$ is called a congruence
 if E is a subalgebra of A .

- In elementary terms, a cong. is an e-rel
 $((x,x) \in E, (x,y) \in E \Rightarrow (y,x) \in E, (x,y) \in E, (y,z) \in E \Rightarrow (x,z) \in E)$
 such that

$s \in \Omega_n, (x_1, y_1) \in E, \dots, (x_n, y_n) \in E,$
 $(s(x_1, \dots, x_n), s(y_1, \dots, y_n)) \in E$.

- I will write $x E y$ to mean $(x, y) \in E$.

• If $E \xrightarrow{i} A^2$ is a congruence, can form diagram

$$E \xrightarrow{i} A^2 \begin{array}{c} \xrightarrow{P_1} \\ \xrightarrow{P_2} \end{array} A \text{ in } \Omega\text{-Alg}$$

domain

& so obtain

$$E \begin{array}{c} \xrightarrow{d} \\ \xrightarrow{c} \end{array} A: (x,y) \in E \begin{array}{c} \xrightarrow{d} x \\ \xrightarrow{c} y \end{array}$$

codomain

a pair of Ω -algebra homomorphisms.

• We are interested in their coequalisers

$A \xrightarrow{p} A/E$ in $\Omega\text{-Alg}$
defined as follows:

- elements of A/E are equiv. classes $[a]$
with $p(a) = [a] = \{x : x E a\}$.

• Observe p is surjective. Therefore if p is to be a homomorphism we are forced to define

$$S^{A/E}([a_1], \dots, [a_n]) = [S^A(a_1, \dots, a_n)].$$

• Is this well defined?

Suppose $[b_1] = [a_1], \dots, [b_n] = [a_n]$
Then $b_1 E a_1, \dots, b_n E a_n$ so as E
a congruence we have

$$s(b_1, \dots, b_n) E s(a_1, \dots, a_n)$$

$$\text{so } [s(b_1, \dots, b_n)] = [s(a_1, \dots, a_n)]$$

as required.

• In particular, A/E is a Ω -algebra

& $p: A \longrightarrow A/E$ a surjective
homomorphism.

Proposition $E \begin{matrix} \xrightarrow{d} \\ \xrightarrow{c} \end{matrix} A \xrightarrow{p} A/E$ is a coequaliser in $\mathcal{U}\text{-Alg}$.

Proof

- Firstly if $(x, y) \in E$ then

$$pd(x, y) = [x] = [y] = pc(x, y)$$

$$\text{so } pd = pc.$$

- Given $A \xrightarrow{f} B$ with $fd = fc$.

This means precisely that if

$$(x, y) \in E \text{ then } fx = fy$$

$$\text{Therefore } [x] = [y] \implies fx = fy.$$

- Therefore we can extend f along p

$$\text{by } \begin{array}{ccc} A & \xrightarrow{p} & A/E \\ & \searrow f & \downarrow \bar{f} \\ & & B \end{array}$$

$$\text{where } \bar{f}[a] = fa$$

- Clearly \bar{f} is a homomorphism, since f is.

- Since p is surjective, \bar{f} is only map extending f along p .

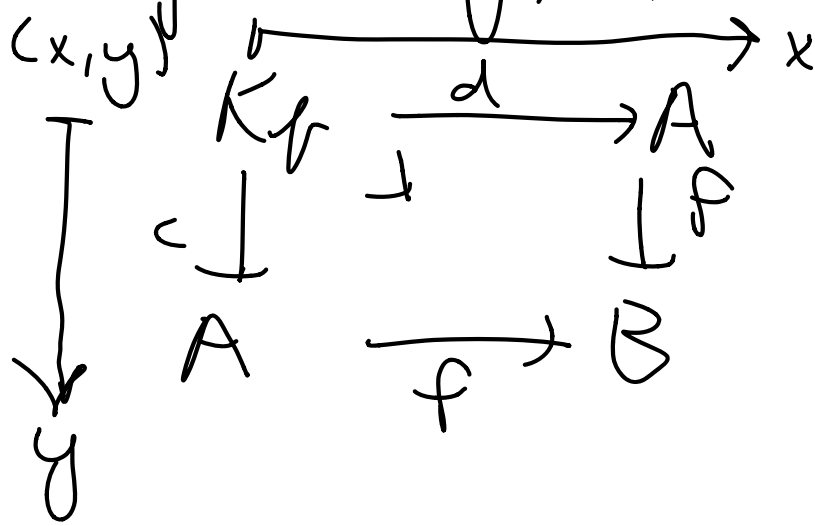
Def.) Let $f: A \rightarrow B \in \mathcal{U}\text{-Alg}$.

The kernel of f is the congruence

$$K_f = \{ (x, y) : fx = fy \} \longleftarrow \downarrow \longrightarrow A^2.$$

- Easy to see this is a congruence :
check it!

Categorically, K_f is the pullback



so, in particular,
we have
 $fd = fc$.

Therefore, we get a unique factorisation of f through the coequaliser



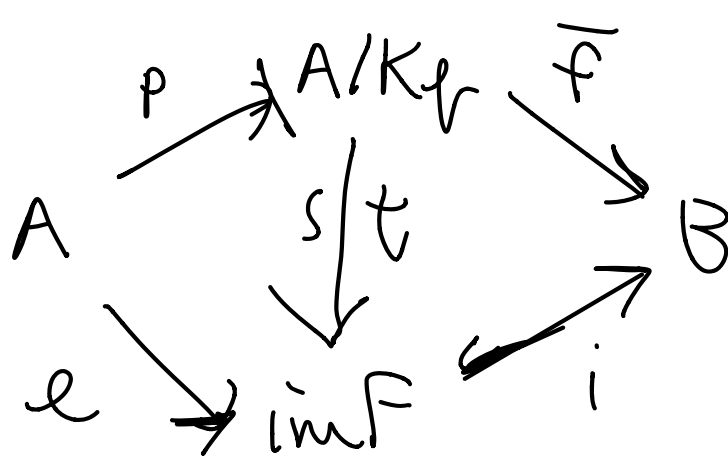
First isomorphism theorem

• Given $f: A \rightarrow B \in \Omega\text{-Alg}$,
we have an isomorphism

$$t: A/K_f \longrightarrow \text{im } f$$

commuting with the factorisations
of $f \cup$: that is,

the diag:



Proof

- Define $t[a] = fa$.

Clearly makes diagram commute.
It is a homom. as f is.

- For surj., if $b \in \text{im } f$, then $b = fa$ so $t[a] = fa = b$.

- For inj., suppose $t[a] = t[b]$.

That is, $fa = fb$.

Then $(a, b) \in K_f$ so $[a] = [b]$.

□

Corollary

If $f: A \rightarrow B$ is surjective,

then $A/K_f \cong B$. In particular,

$K_f \xrightarrow{d} A \xrightarrow{f} B$ is a coequaliser

Proof

In this case $\text{im } f = B$

$$\text{so } A/K_f \xrightarrow{\sim} B$$

$$\begin{array}{ccc} & p & \\ & \searrow & \\ \text{so } A & & A/K_f \\ & f & \downarrow \text{st} \\ & & B \end{array}$$

& then use that coequalisers are invariant up to iso, so as p coequaliser map, $f = \underset{\text{iso}}{t \circ p}$ is also coequaliser.