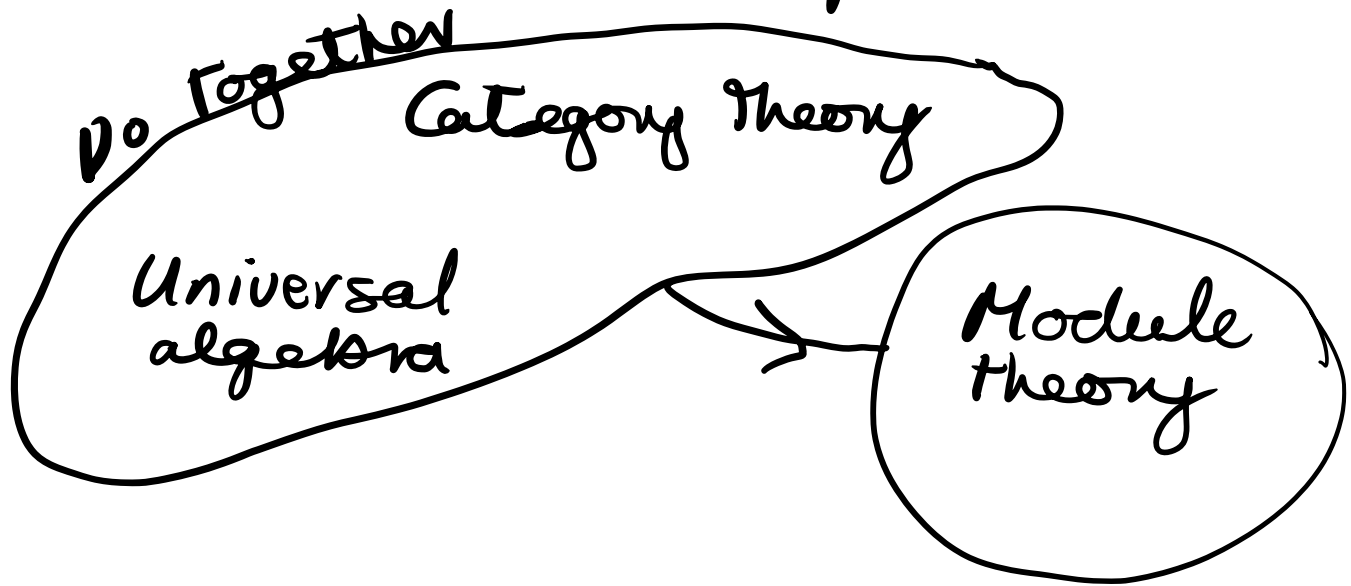


Algebra 3

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- lecture notes + IS each week

- Course has 3 components:

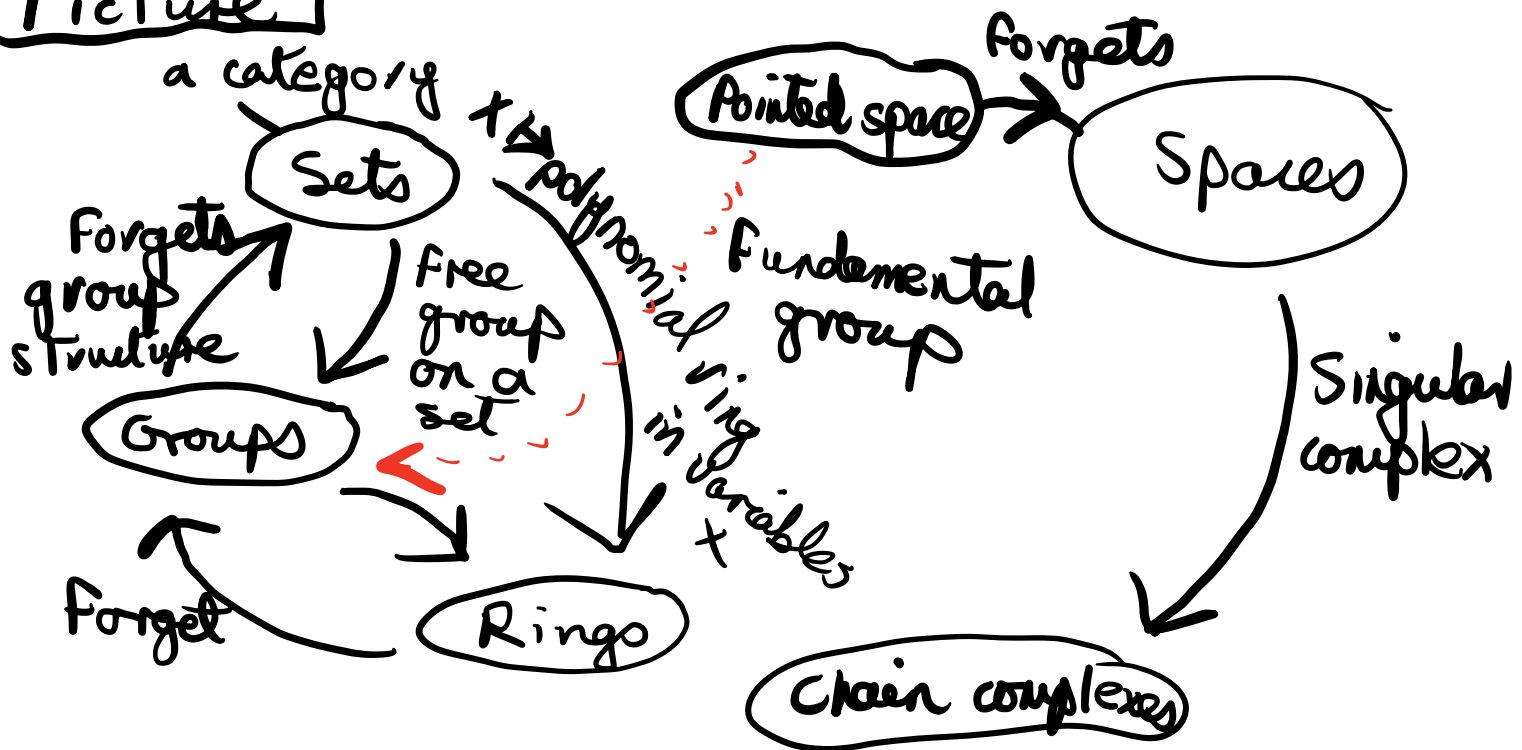


Today : start with category theory .

What is category theory?

- In math, study structures like sets, groups, rings, topological spaces.
- Each forms category of "structures of that type".
- Cat. Theory studies the relⁿ between these diff. areas of mathematics (or categories of mathematics)

Picture



- What do these different categories have in common?
- What structure is preserved when we move between them?
- Fundamental notion in category theory is an arrow/morphism:
$$A \longrightarrow B$$

between two things ~ captures relationship.

Def) A category $\mathcal{C}$ consists of :

- a collection $\text{ob } \mathcal{C}$ of objects $A, B, C, \dots$
- & for each pair of objects $A, B \in \text{ob } \mathcal{C}$ a collection $\mathcal{C}(A, B)$ of "arrow/morphisms" from A to B ,
(depicted $A \xrightarrow{F} B$ or $F: A \rightarrow B$ to mean $F \in \mathcal{C}(A, B)$)

- For each $A, B, C \in \text{ob } \mathcal{C}$ a function $\mathcal{C}(B, C) \times \mathcal{C}(A, B) \longrightarrow \mathcal{C}(A, C)$

$(B \xrightarrow{g} C, A \xrightarrow{f} B) \longmapsto A \xrightarrow{g \circ f} C$
called composition.

- For each object $A \in \text{ob } \mathcal{C}$ an arrow $1_A: A \rightarrow A$ called the identity on A .

- These satisfy the following axioms:
Associativity

- Given $A \xrightarrow{f} B, B \xrightarrow{g} C, C \xrightarrow{h} D$
we have $(h \circ g) \circ f = h \circ (g \circ f)$.

- Left & right unit laws

- Given $f: A \rightarrow B$ we have
 $1_B \circ f = f = f \circ 1_A$.

- Notation: Often write $gf: A \rightarrow C$ instead of $g \circ f$.
- Associativity & unit laws imply that given

$$A_0 \xrightarrow{f_1} A_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} A_n$$

there is a unique way to compose the arrows $f_n f_{n-1} \dots f_2 f_1$, indep. of brackets & identities.

Eg: When $n=4$,

$$((f_4 f_3) f_2) f_1 = (f_4 (A_3)) ((f_3 f_2) f_1).$$

- Commutative diagrams

Given a diagram such as

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ h \downarrow & & \downarrow g \\ C & \xrightarrow{i} D \xrightarrow{j} & E \end{array}$$

we say the diagram commutes

if the paths agree:

in this case, $gf = jih$.

- In a larger diagram, eg.

$$\begin{array}{ccccc} A & \xrightarrow{f} & B & \xrightarrow{k} & F \\ h \downarrow & \parallel & \downarrow g & \parallel & \downarrow l \\ C & \xrightarrow{i} & D \xrightarrow{j} & E & \xrightarrow{m} G \end{array}$$

diag. commutes

The commutativity of subdiagrams implies the commutativity of the outer diagram:

This is diagram chasing.

Examples

- Set : objects are sets,
morphisms $f: A \rightarrow B$ are
Functions.
- Given $A \xrightarrow{f} B$ & $B \xrightarrow{g} C$ the function
 $gf: A \rightarrow C$ is def. by $gf(x) = g(f(x))$.
We have $1_A(x) = x$.

Categories of sets with structure

- Mon, the cat of monoids &
monoid homomorphisms:
 $f: (A, m_A, e_A) \rightarrow (B, m_B, e_B)$
 $f(e_A) = e_B$
 $f(m_A(x, y)) = m_B(fx, fy)$
- Grp, the category of groups
& group homomorphisms.
- Rng ~ rings, ring hom...
- $K\text{-Vect}$ ~ K -vector spaces, lin. transf.
- These are examples of
algebraic categories.

- More generally, given signature (Σ, E) set of equations we can consider the cat. $(\Sigma, E)\text{-Alg}$. This captures all of the above examples. Return to this example later.

- Top is the category of topological spaces and continuous functions.

Def) A morphism $f: A \rightarrow B \in \mathcal{C}$
 is an isomorphism if
 $\exists g: B \rightarrow A$ such that
 $gf = 1_A$ & $fg = 1_B$.

Remark) Can express this via the
 diagram

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ & \searrow 1_A & \downarrow g \\ & & A \xrightarrow{f} B \end{array}$$

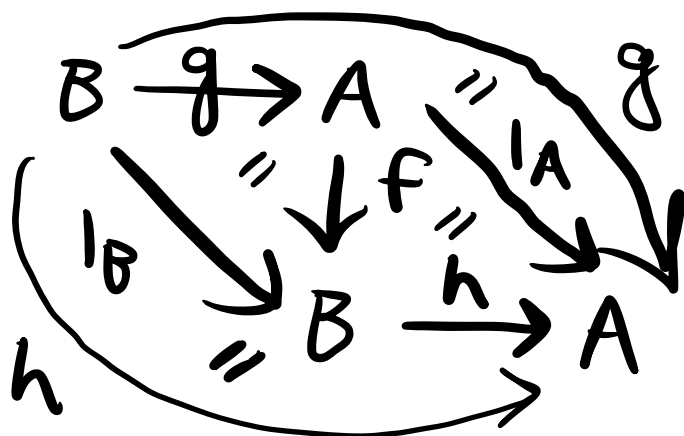
(Note: In the original image, there are additional arrows and labels: a diagonal arrow from A to B labeled 1_B, a diagonal arrow from B to A labeled g, and a diagonal arrow from A to B labeled f. The diagram is a commutative square with A at top-left, B at top-right, A at bottom-left, and B at bottom-right. Arrows are: A to B (f), A to A (1_A), B to B (1_B), and B to A (g).)

We say that g is the inverse of
 f and write $g = f^{-1}$.

This is justified by:

Proposition) If $f: A \rightarrow B$ is an iso,
 then its inverse is unique.

Proof) Suppose $g, h: B \rightarrow A$
 are inverses to f .



Therefore
 $g = h$.

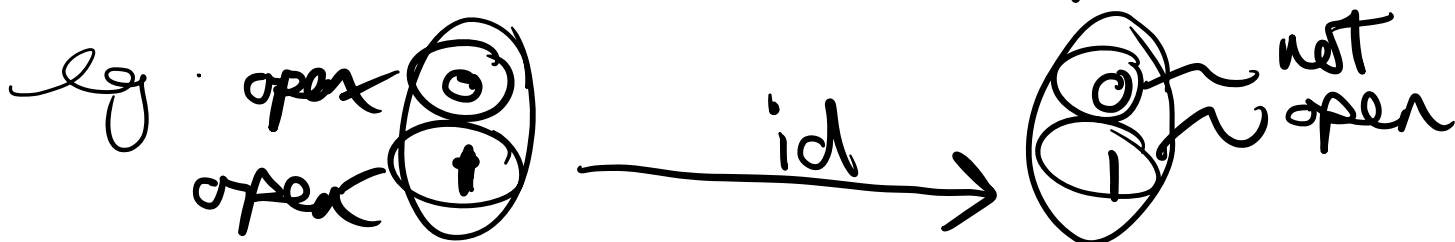
In alg. notation, the corresponding proof is

$$g = 1_A \circ g = (h \circ f) \circ g = h \circ (f \circ g) = h \circ 1_B = h.$$

Examples

- In Set, the isomorphisms $f: A \rightarrow B$ are the bijections.
- In other algebraic cats such as Grp, an isomorphism is sometimes defined as a bijective homomorphism: this coincides with the categorival notion, since if a homomorphism is bijective its inverse (as a function) is also a homomorphism.
- For Top, an isomorphism is a homeomorphism: a cts function with a cts inverse.

In Top, $\exists$ bijective cts maps which are not homeomorphisms



not a homeomorphism.

- In summary, the categorical defⁿ of isomorphism captures the correct notion in all of our examples.

Defⁿ) - A category $\mathcal{C}$ is said to be locally small if for all $A, B \in \mathcal{C}$ the collection $\mathcal{C}(A, B)$ is a set.

- If, furthermore, the collⁿ of $\mathcal{C}$ is a set, we say that it is small.

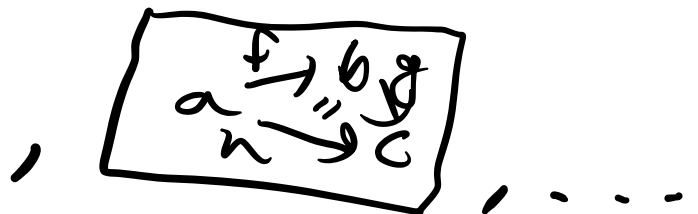
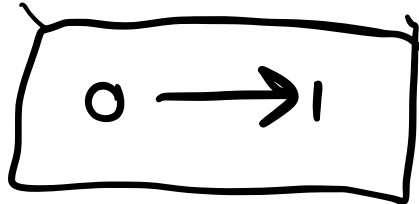
Examples of locally small cats

- Set is not small: one cannot form a "set of all sets". However, it is locally small as the collection $\text{Set}(A, B)$ of functions from A to B forms a set.
- All of the examples considered so far are locally small, but not small.

Examples of small categories

Ex 1) - $\emptyset$ empty cat

- $I = (\bullet)$ 1 object, 1 identity morph
- plus identities



Ex. 2

Preorder & posets

Preorder $(X, \leq)$
satisfying $x \leq x$

$$x \leq y \text{ \& } y \leq z \Rightarrow x \leq z$$

Poset: also $x \leq y \text{ \& } y \leq x \Rightarrow x = y$.

- If $(X, \leq)$ a preorder,
can form a category X^*
whose objects are elements of X
& such that there exists a single
morph. $x \rightarrow y \iff x \leq y$ &
no morphisms from x to y
otherwise.

- In fact,
Preorders $\overset{\text{same}}{\equiv}$
hopefully later in course
(equivalence cats)

Small cats with
at most 1 morph.
between any
two objects.

Example 3 ^{mult, unit}

If $(M, \cdot, e)$ is a monoid
we can form a category ΣM w'
one object $\bullet$ & whose
morphisms $\bullet \xrightarrow{m} \bullet$ are the
elements of M .

We compose by $\bullet \xrightarrow{m} \bullet \xrightarrow{n} \bullet \xrightarrow{e} \bullet$
unit is $1_\bullet$

& assoc. & unit laws for monoid
 $\Rightarrow$ axioms for a category

- If $(M, \cdot, e)$ is a group (all elts. have inverse)
then ΣM is a groupoid: a category
in which all morphisms
are isomorphisms.

- Monoids $\equiv$ 1-object small cats
Groups $\equiv$ 1-ob. small groupoids

- Natural: e.g. symmetric group S_n
consists of the bijections

$$\bar{n} = \{1, \dots, n\}$$

$\curvearrowright$ Groups of transformations
 $\sim$ symmetries of an object.

Functors

Def") Let A, B be categories. A Functor $F: A \longrightarrow B$ consists of:

- a function $\text{ob } A \longrightarrow \text{ob } B$ which we write as $x \longmapsto Fx$,
- For each pair $x, y \in A$ a function $A(x, y) \longrightarrow B(Fx, Fy)$ written as $x \xrightarrow{\alpha} y \longmapsto Fx \xrightarrow{F\alpha} Fy$.

This satisfies the axioms:

- given $x \xrightarrow{\alpha} y \xrightarrow{\beta} z \in A$, we have $F(\beta \circ \alpha) = F\beta \circ F\alpha$.
- given $x \in A$ we have $F(1_x) = 1_{Fx}$.

In other words, a functor sends obs to obs, arrows to arrows (respecting their domain & codomain) & preserves identities and composition.

Examples

- Forgetful Functors:

Eg. $U: \text{Grp} \longrightarrow \text{Set}$
 $(X, m, e) \longmapsto X$
mult. unit

which sends a group to its underlying set, i.e. forgets the group structure.

Sim. it sends a group homomorphism to its underlying function:

$$f: (X, m_X, e_X) \longrightarrow (Y, m_Y, e_Y) \longmapsto f: X \longrightarrow Y$$

Clearly preserves composition & identities.

- Similarly, there are forgetful functors from cats of "sets with structure"

To Set, eg $U: \text{Top} \rightarrow \text{Set}$, $U: \mathbf{K}\text{-Vect} \rightarrow \text{Set}$
 or $\mathbf{Rng}$, $\mathbf{Mon}$ etc...
 Similarly, $U: \mathbf{Rng} \rightarrow \mathbf{Mon}$
 $(\mathbf{R}, +, \times, 0, 1) \mapsto (\mathbf{R}, \times, 1)$.

There is a functor $F: \text{Set} \rightarrow \mathbf{Mon}$

where $FX =$ "list/free" monoid on X

- Its elements are (possibly empty) lists $[x_1, x_2, \dots, x_n]$ of elements of X with unit $[-]$ empty list.
- Composition of lists is "concatenation":
 $[x_1, \dots, x_n] \cdot [y_1, \dots, y_m] = [x_1, \dots, x_n, y_1, \dots, y_m]$
- This makes FX a monoid.
- Given a function $f: X \rightarrow Y$ obtain a monoid homomorphism
 $ff: FX \rightarrow FY$

$$[x_1, \dots, x_n] \mapsto [fx_1, \dots, fx_n]$$

& easy to check that this satisfies axioms for a Functor.

Example) Let G be a group & consider its suspension ΣG as a groupoid with one object

• & $\Sigma G(-, -) = G$.

Then a functor $\Sigma G \rightarrow \text{Set}$ is what?

$$g \downarrow \mapsto \begin{array}{c} X \\ \downarrow g \cdot - \\ X \end{array} \quad \begin{array}{c} x \\ \downarrow \\ g \cdot x \end{array}$$

The functor axioms

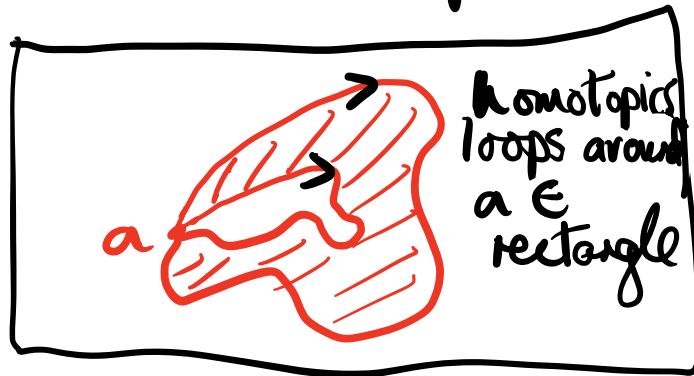
say $e \cdot x = x$ & $h.(g \cdot x) = (h.g) \cdot x$.

Therefore a functor $\Sigma G \longrightarrow \text{Set}$
 is exactly a G-set: a set X with
 an action of the group G .
 What is a functor $\Sigma G \longrightarrow \text{Vect}$?

Example Let Top_* be category of
 pointed spaces (X, a)
 $\begin{matrix} \text{Top space} \\ \swarrow \\ X \end{matrix}$ $\begin{matrix} a \in X \\ \searrow \\ \text{basepoint} \end{matrix}$
 & continuous maps which
 preserve basepoint.

There is a functor $\Pi_1 : \text{Top}_* \longrightarrow \text{Grp}$
 where $\Pi_1(X, a) =$ fundamental
 group of
 loops around a (up to homotopy)

eg.



$$\Pi_1(S^1, *) = \mathbb{Z}$$

There are also functors
 $H_n : \text{Top} \longrightarrow \text{Ab} \sim \text{abelian group}$
 sending a top. space
 to its n 'th homological group.

Example : CAT

- Given functors $A \xrightarrow{F} B \xrightarrow{G} C$
we can compose them in obvious way :


$$x \mapsto G(F(x)) = GFX$$

$$x \xrightarrow{\alpha} y \mapsto GFX \xrightarrow{G\alpha} GFy$$

- Likewise, we have identity functor
 $1_A: A \rightarrow A : x \mapsto x$
- In this way, we obtain a (large) category CAT of categories & functors. In fact, this is a 2-category!